

NATIONAL BUREAU OF STANDARDS REPORT

8171

A PROGRAMMED FORMALIZER FOR A FRAGMENT OF ENGLISH

by

Sylvan Cappell

To the

National Science Foundation
(Grant No. GN 107)



U. S. DEPARTMENT OF COMMERCE
NATIONAL BUREAU OF STANDARDS

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NBS PROJECT

11-11-11404

January 15, 1964

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U. S. DEPARTMENT OF COMMERCE
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A PROGRAMMED FORMALIZER FOR A FRAGMENT OF ENGLISH

This paper describes a computer program written to translate English into symbolic logic by an algorithm developed by Walter Sillars.^{1/} This algorithm, termed "formalizer," has been written for use in a picture language machine and the vocabulary it applies to is directed to that purpose.^{2/} As input the program accepts a parsed sentence, one which has been analyzed grammatically. The parser used is described in a paper by Donald Cohen.^{3/} The formalizer program is quite general and, like the parser, can be modified or expanded for other grammars than the one discussed here.

Sillars' paper describes a formalizer written for a grammar for a fragment of English, Grammar 12R. A corrected version of the grammar and formalizer appears in Appendix I. Grammar 12R is a subset of Grammar 12 developed by B. Kirk Rankin, III.^{4/}

The formalizer program was written in the COMIT programming language, knowledge of which is assumed in this paper. The program appears in Appendix II together with model input and output.

The parsed input describes which rules of Grammar 12R would be used in the generation of the sentence used as input to the parser. The existence of discontinuous rules in a grammar would necessitate the rearrangement of the rule numbers in the output of the parser, in the general case. However, for the special case of Grammar 12R, this is unnecessary.

The formalizer program will accept any number of parser inputs which are separated by a card with the word NEW on it. The program terminates when there are no more input cards to be read in.

The output of the formalizer program is in symbolic logic. A variable is symbolized by a subscripted X. UQ is the symbol used to represent the universal quantifier, EQ the existential quantifier, UN the Sillars U quantifier (Russell and Whitehead's E!), CNJ the "conjunction" and * - represents Church's T. All other symbols are the same as Sillars'.

Throughout most of the program, what Sillars calls pseudo-wffs are stored in the workspace. The program uses only shelves 1-6.

The first part of the program, the part before the rule named GRAM, reads in the input cards until it reaches one with NEW on it. The program assumes that consecutive digits form part of the same rule number and that non-digital characters separate any two rule numbers. The program assembles the rule numbers in the order that they are to be used on shelf 1 and places an *B after the last rule number. An X_1 (x/.1 in COMIT) is stored on shelf 3. Whenever Sillars' algorithm calls for a new variable the contents of shelf 3 are used and the subscript value of the X in shelf 3 is increased by 1. The program places the initial symbol, TOP, inclosed in brackets in the workspace. Brackets inclose all pseudo-wffs. *B is used internally to represent [and *C represents] . In the output, however, they are converted to "(" and ")" respectively.

The program then branches to the rule named EXIT. It picks up the first grammar rule number on shelf 1 and looks it up in the list GRAM. It then branches to the formalizer instructions corresponding to the grammar rule number, erases the rule number and returns to EXIT. Those grammar rule numbers which have no corresponding formalizer instructions, the ones which Sillars calls not applicable, NA, are simply erased and control returns to EXIT.

After the last grammar rule number has been picked up the program picks up the *B on shelf 1 and finds it in the list GRAM. It then branches to STOP and the symbolic logic statement corresponding to the original sentence is given as output. The program then prepares to read in the next input sentence.

The subroutine named FUNC (the name of its first rule) locates the shortest pseudo-wff containing all occurrences of NV and places an *M immediately before it and an *N immediately after it. It terminates one of the subrules of OUT, which then branches back into the main body of the program.

The program can be modified quite readily. For each new rule assign a new number. Place a card under GRAM with the number on the left and a zero in the right half and branch to a point in the program where the actual formalizer rules are to be executed. Control must finally return to EXIT.

If an instruction requires finding the shortest pseudo-wff containing all occurrences of some constituent α , substitute NV for

each α , indicate a subrule of OUT to be used, and branch to FUNC.
(Subrules can, of course, be added to OUT.) If so desired, every NV
can then be replaced by α .

To substitute an X_1 for each NV and to then increase the subscript
of the X in shelf 3 by 1, branch to XOG. Control will automatically be
returned to EXIT.

The program treats any symbol beginning with PC or SC as if it
consisted only of one of those two respective symbols. The processing
of a PC rule without a quantifier begins at APC1 and with a quantifier
at APC2 and all SC rules at ASC. Each rule number associated with a
DEP or DES rule leads in the program to a series of instructions where
quantifiers are generated. The rules UNAT, SQUAT, TQUAT and NQUAT
substitute the appropriate quantifier and branch to QUADS where CNJ
is substituted for \square . UQUAT substitutes a UQ for Q and PLY for \square .
More quantifier rules, accomplishing similar functions, can be added
to the program.

A rule can be taken out simply by removing the card with the
appropriate rule number under GRAM.

In the final analysis this program is to be used in the predicate
evaluator in the picture language machine. The output format in which
all propositions which form part of other propositions are inclosed in
parentheses should facilitate its use for this application. Thus a
proposition of the general form $p \wedge q \supset r$ would be given in output as
 $((p) \wedge (q)) \supset r$.

In determining the truth or falsehood of a proposition which has several component propositions, the predicate evaluator must selectively choose which of the latter to begin with. Thus in evaluating propositions of the form $p \vee q \supset r$ it will usually be advisable to begin by analyzing the truth-value of r . If, however, r is in some sense very complex, it might be preferable to begin with p .

APPENDIX I. The Grammar and Formalizer

Here is the revised grammar for which the program in Appendix II has been written. For the notational conventions, see Sillars.^{1/}

TOP	=	CL1	
TOP	=	CL3	
TOP	=	CL5	
TOP	=	CL7	
TOP	=	CL11	SR
TOP	=	CL13	
TOP	=	CL15	
TOP	=	CL17	
AAP1	=	WPN + COM1	AAP1 = WNP + COM1
AAP1	=	COM1	SR
AAP1	=	WOP + J	AAP1 = WOP + J
AAS	=	WN + COM1	AAS = WN + COM1
AAS	=	COM1	SR
AAS	=	WO + J	AAS = WO + J
APN	=	ET + PNBET	SR(2)
AR	=	are	NA
ARNTC	=	aren't	

$$\begin{array}{rcl}
 \text{CL1} & = & \text{SUB1X} + \text{PRE0} \\
 \text{CL3} & = & \text{SUB6X} + \text{PRE1} \\
 \text{CL5} & = & \text{SUB4X} + \text{PRE3} \\
 \text{CL7} & = & \text{SUB7W} + \text{PRE4} \\
 \text{CL11} & = & \text{SUB10} + \text{PRE11} \\
 \text{CL13} & = & \text{SUB11} + \text{PRE13} \\
 \text{CL15} & = & \text{SUB13} + \text{PRE15} \\
 \text{CL17} & = & \text{SUB16} + \text{PRE17}
 \end{array}
 \left. \vphantom{\begin{array}{rcl} \text{CL1} \\ \text{CL3} \\ \text{CL5} \\ \text{CL7} \\ \text{CL11} \\ \text{CL13} \\ \text{CL15} \\ \text{CL17} \end{array}} \right\} \text{CL}_i = \text{PRE}_j(\text{SUB}_k)$$

$$\text{CO} = , \quad \text{NA}$$

$$\text{COL} = \text{Color} \quad \text{COL} = \text{Smc}$$

$$\text{COL} = \text{Size} \quad \text{COL} = \text{Smz}$$

$$\begin{array}{rcl}
 \text{COM1} & = & \text{LP} \\
 \text{COM1} & = & \text{RE} \\
 \text{COM2} & = & \text{COM1} \\
 \text{COM2} & = & \text{PNP2} \\
 \text{COM3} & = & \text{COM1} \\
 \text{COM3} & = & \text{SNP2}
 \end{array}
 \left. \vphantom{\begin{array}{rcl} \text{COM1} \\ \text{COM1} \\ \text{COM2} \\ \text{COM2} \\ \text{COM3} \\ \text{COM3} \end{array}} \right\} \text{SR}$$

DEP12 = some
 DEP12 = the
 DEP12 = two
 DEP12 = all
 DEP12 = no
 DEP2 = some
 DEP2 = the
 DEP2 = two
 DEP2A = any
 DEP5S = some
 DEP5S = two
 DEP5S = no
 DES1 = a
 DES1 = one
 DES13 = the
 DES13 = a
 DES13 = every
 DES13 = one
 DES13 = each
 DES16 = the
 DES16 = a
 DES16 = one

Procedure for rewriting (QNV) [NN(NV)

□ A(NV)] where NN is either N or NPLUR:

DE is any DEPi or DESi

1. If DE = the, Q = U

If DE = No or any, (QNV) := T (ENV)

If quantifier is deleted, Q = E. In

all other cases,

if DE =	then Q =
---------	----------

some	S
------	---

the	S
-----	---

two	T
-----	---

all	V
-----	---

a	E
---	---

one	U
-----	---

each	V
------	---

every	V
-------	---

Then replace all instances of NV by the next (in alphabetic order) individual variable which has not yet been used.

2. If DE = all, each or every, then

□ = ∅., otherwise □ = &.

3. Replace every expression of the form Vi(Vj), where Vi and Vj are individual variables, by (Vi = Vj).

DES16	=	each	
DES16	=	every	
DES16	=	no	
DES2	=	the	
DES2	=	a	
DES2	=	one	
DES7	=	a	
DES7	=	one	
DES7	=	no	
DUM1	=	R1 + NPH	DUM1(a,NPH) = R1(a, NPH1, NPH2)
DUM2	=	SNASN	SR
DUM2	=	PNAPN	
DUM2	=	PNBET	
ES	=	ET + SNP2	SR(2)
ET	=	and	NA
F1	=	T2 + G1	SR(2)
F2	=	T2 + G2	
G1	=	right	MT + G1 = Mort; G1 = Rt
G1	=	left	MT + G1 = Molf; G1 = Lf
G1	=	top	MT + G1 = Mtop; G1 = Tp
G1	=	bottom	MT + G1 = Mbot; G1 = Bot
G2	=	center	MT + G2 = Mcen; G2 = Cen
G2	=	middle	MT + G2 = Mmid; G2 = Mid
I	=	is	NA

IA	=	J1	}	SR
IA	=	J2		
IA	=	J1 + J2	IA(a) = J1(a) & J2 (a)	
ISNTC	=	isn't	NA	
J	=	J1	}	SR
J	=	J2		
J1	=	big	J1 = Bg	
J1	=	little	J1 = Lt	
J1	=	large	J1 = Lg	
J1	=	small	J1 = Sm	
J2	=	black	J2 = Bk	
J2	=	white	J2 = Wh	
JER	=	bigger	JER= Bgr	
JER	=	littler	JER= Ltr	
JER	=	larger	JER= Lgr	
JER	=	smaller	JER= Smr	
LL	=	Z1	}	SR
LL	=	Z2		
LP	=	LL		
MO	=	more	NA	
MT	=	MO : TN	NA	
N	=	triangle	N = Tr	
N	=	square	N = Sq	
N	=	circle	N = Cir	

NPH	=	SNP2	SR
NPH	=	PNP2	SR
NPLUR	=	triangles	NPLUR = Tr
NPLUR	=	squares	NPLUR = Sq
NPLUR	=	circles	NPLUR = Cir
NT	=	not	NT = \neg
P	=	in	P = In
P	=	near	P = Nr
P	=	below	P = Bel
P	=	above	P = Ab
P	=	touching	P = Tch
P1	=	to	$\left. \begin{array}{l} \\ \\ \\ \end{array} \right\} \text{NA}$
P1	=	on	
P1	=	at	
P2	=	in	
P3	=	between	P3 = Bet
P4	=	of	NA
PA1	=	IA : AAP1	PA1(a) = IA(a)&AAP1(a)
PA1	=	\emptyset : AAP1	$\left. \begin{array}{l} \\ \end{array} \right\} \text{SR}$
PA1	=	IA	
PA2	=	IA : AAP1	PA2(a) = IA(a)&AAP1(a)
PA2	=	IA	SR

PC12	=	DEP12:NPLUR	} A(PCi) = (QNV) [NPLUR(NV) □ A(NV)]
PC2	=	DEP2 :NPLUR	
PC2	=	∅ :NPLUR	
PC2A	=	DEP2A : NPLUR	
PC2A	=	∅ : NPLUR	
PC4	=	DEP2 + NPLUR	
PC4	=	NPLUR	
PC412	=	DEP12 + NPLUR	
PC4A	=	DEP2A + NPLUR	
PC4A	=	NPLUR	
PC4S5	=	DEP5S + NPLUR	
PC4S5	=	NPLUR	
PC4S6	=	PC4S5	SR
PNAPN	=	PNBET + APN	A(PNAPN) = A(PNBET, APN)
PNBET	=	PC2 + PA1	A(PNBET) = PA1(PC2) & A(PC2)
PNBET	=	PC4	SR
PNP12	=	PC412	SR
PNP12	=	PC12 + PA1	A(PNP12) = PA1(PC12) & A(PC12)
PNP2	=	PC2 + PA1	A(PNP2) = PA1(PC2) & A(PC2)
PNP2	=	PC4	SR
PNP2A	=	PC2A + PA1	A(PNP2A) = PA1(PC2A) & A(PC2A)
PNP2A	=	PC4A	SR
PNP6S	=	PC4S6	SR

PPLB	=	P3 + DUM2	PPLB(a) = P3(a, DUM2)
PREO	=	VPO	
PRE1	=	VP1	
PRE11	=	VP11	
PRE13	=	VP13	
PRE15	=	VP15	
PRE17	=	VP17	SR
PRE3	=	VP3	
PRE4	=	VP4	
R	=	R1	
R	=	R2	
R	=	R4	
R	=	R5	
R1	=	LL+P4	SR(1)
R2	=	JER + TN	SR(1)
R4	=	P	SR
R5	=	MT + DUM1	R5(a, b) = MT DUM1(a, b)
RE	=	R + NPH	RE(a) = R(a, NPH)
RE	=	PPLB	SR
SA	=	IA : AAS	SA(a) = IA(a) & AAS(a)
SA	=	IA	SR
SA	=	∅ : AAS	SR
SAME	=	same	NA

SC1	=	DES1 : N	}	$A(SC_i) = (QNV) [N(NV) \square A(NV)]$
SC13	=	DES13 : N		
SC16	=	DES16 : N		
SC2	=	DES2 : N		
SC3	=	DES1 + N		
SC37	=	DES7 + N		
SC4	=	DES2 + N		
SC413	=	DES13 + N		
SC416	=	DES16 + N		
SC7	=	DES7 : N		
SNASN	=	SNP2 + ES		$A(SNASN) = A(SNP2, ES)$
SNP1	=	SC1 + SA		$A(SNP1) = SA(SC1) \& A(SC1)$
SNP1	=	SC3		SR
SNP13	=	SC413		SR .
SNP13	=	SC13 + SA		$A(SNP13) = SA(SC13) \& A(SC13)$
SNP16	=	SC416		SR
SNP16	=	SC16 + SA		$A(SNP16) = SA(SC16) \& A(SC16)$
SNP2	=	SC4		SR
SNP7	=	SC37		SR
SNP7	=	SC7 + SA		$A(SNP7) = SA(SC7) \& A(SC7)$

SUB1	=	PNP2A	}	SR
SUB10	=	PNP2		
SUB11	=	PNP12		
SUB13	=	SNP13		
SUB16	=	SNP16		
SUB1X	=	SUB1Y		
SUB1Y	=	THERE : SUB1		SR(2)
SUB3	=	SNP1	}	SR
SUB4X	=	SUB4Y		
SUB4Y	=	THERE : SUB3		SR(2)
SUB6	=	PNP6S		SR
SUB6X	=	SUB6Y		SR
SUB6Y	=	THERE : SUB6		SR(2)
SUB7	=	SNP7		SR
SUB7W	=	SUB7Y		SR
SUB7Y	=	THERE : SUB7		SR(2)
T2	=	the	}	NA
TH	=	that		
TH	=	which		
THERE	=	there		
THSAM	=	T2 + SAME	}	NA
TN	=	than		

VPO	=	ARNTC : COM1	VPO = TCOM1
VP1	=	AR : COM1	SR(2)
VP11	=	ARNTC + COM2	VP11 = T COM2
VP13	=	AR + COM2	SR(2)
VP15	=	ISNTC + COM3	VP15 = T COM3
VP17	=	I + COM3	SR(2)
VP3	=	ISNTC : COM1	VP3 = T COM3
VP4	=	I : COM1	SR(2)
WA	=	TH + AR	NA
WI	=	TH + I	NA
WN	=	WI + NT	WN = NT
WN	=	WI	WN = \emptyset
WN	=	NT	WN = NT
WNP	=	WA + NT	WNP = NT
WNP	=	WA	WNP = \emptyset
WNP	=	NT	WNP = NT
WO	=	WI + NT	WO = NT
WO	=	WI	WO = \emptyset
WOP	=	WA + NT	WOP = NT
WOP	=	WA	WOP = \emptyset
Z1	=	P1 + F1	SR(2)
Z2	=	P2 + F2	

APPENDIX II. The Formalizer Program

The COMIT program which realizes the algorithm is listed below. To illustrate the output format of the program, we present the formalization of the sentence:

There aren't any triangles on the right.

The parse of this sentence is:

```
2(39(351(352(381(-THERE))(339(221(195(160(-TRIANGLES)))))))
(239(394(35(-AREN*T))(65(148(146(427(172(-ON))(119(378(-THE))
(121(-RIGHT))))))))))
```

(Where the numbers are names of the rules of 12R which were used in generating it.) The output for this sentence from the formalizer program is:

```
*- + *( + *( + EQ + X / .1 + *) + *( + *( + TR + *( + X /
.1 + *) + * + CNJ + *( + RT + *( + X / .1 + *) + *) + *)
+ *) +
```

APPENDIX II

FORMALIZER

STAR	\$=*B+TOP+*C+X/.1		//*Q2 1 2 3,*Q3 4	*
BEGIN	\$=		//*RCR1	BEGINA
*				FINISH
BEGINA	N+E+W=1		//*A4 1	SUB
BEGINB	\$		//*Q4 1	BEGIN
SUR	\$1		//*L1	ZAS
*	\$=*B		//*Q1 1,*A2 1	*
EXIT	\$=*A+1		//*N1 1,*L 1	GRAM
-ZAS	*0		//*Q4 1	SUB
	*1		//*Q4 1	SUB
	*2		//*Q4 1	SUB
	*3		//*Q4 1	SUB
	*4		//*Q4 1	SUB
	*5		//*Q4 1	SUB
	*6		//*Q4 1	SUB
	*7		//*Q4 1	SUB
	*8		//*Q4 1	SUB
	*9		//*Q4 1	SUB
*	\$1=		//*A4 1,*K1,*Q1 1	SUB
-GRAM	*2	=0		R2
	*4	=0		R4
	*6	=0		R6
	*8	=0		R8
	*1*0	=0		R10
	*1*2	=0		R12
	*1*4	=0		R14
	*1*6	=0		R16
	*2*5	=0		R25
	*2*6	=0		R26
	*2*7	=0		R27
	*2*8	=0		R28
	*2*9	=0		R29
	*3*0	=0		R30
	*3*3	=0		R33
	*3*9	=0		R39
	*5*3	=0		R53
	*5*9	=0		R59
	*6*1	=0		R61
	*4*1	=0		R41
	*4*3	=0		R43
	*4*5	=0		R45
	*4*7	=0		R47
	*6*3	=0		R63
	*6*4	=0		R64
	*6*5	=0		R65
	*6*6	=0		R66
	*6*7	=0		R67
	*6*8	=0		R68
	*6*9	=0		R69
	*7*0	=0		R70
	*7*1	=0		SQUAT
	*7*2	=0		SQUAT
	*7*3	=0		TQUAT
	*7*4	=0		UQUAT

*7*5	=0	NQUAT
*7*6	=0	SQUAT
*7*7	=0	SQUAT
*7*8	=0	TQUAT
*8*3	=0	NQUAT
*8*9	=0	SQUAT
*9*0	=0	TQUAT
*9*1	=0	NQUAT
*9*2	=0	EQUAT
*9*3	=0	UNAT
*9*4	=0	SQUAT
*9*5	=0	EQUAT
*9*6	=0	UQUAT
*9*7	=0	UNAT
*9*8	=0	UQUAT
*9*9	=0	SQUAT
*1*0*0	=0	EQUAT
*1*0*1	=0	UNAT
*1*0*2	=0	UQUAT
*1*0*3	=0	UQUAT
*1*0*4	=0	NQUAT
*1*0*5	=0	SQUAT
*1*0*6	=0	EQUAT
*1*0*7	=0	UNAT
*1*0*8	=0	EQUAT
*1*0*9	=0	UNAT
*1*1*0	=0	NQUAT
*1*1*1	=0	R111
*1*1*2	=0	R112
*1*1*3	=0	R113
*1*1*4	=0	R114
*1*1*7	=0	R117
*1*1*9	=0	R119
*1*2*0	=0	R120
*1*2*1	=0	R121
*1*2*2	=0	R122
*1*2*3	=0	R123
*1*2*4	=0	R124
*1*2*5	=0	R125
*1*2*6	=0	R126
*1*3*0	=0	R130
*1*3*1	=0	R131
*1*3*2	=0	R132
*1*3*4	=0	R134
*1*3*5	=0	R135
*1*3*6	=0	R136
*1*3*7	=0	R137
*1*3*8	=0	R138
*1*3*9	=0	R139
*1*4*0	=0	R140
*1*4*1	=0	R141
*1*4*2	=0	R142
*1*4*3	=0	R143
*1*4*4	=0	R144
*1*4*5	=0	R145

*1*4*6	=0	R146
*1*4*7	=0	R147
*1*4*8	=0	R148
*1*5*2	=0	R152
*1*5*3	=0	R153
*1*5*4	=0	R154
*1*5*6	=0	R156
*1*5*8	=0	R158
*1*5*9	=0	R159
*1*6*0	=0	R160
*1*6*1	=0	R161
*1*6*2	=0	R162
*1*6*4	=0	R164
*1*6*6	=0	R166
*1*6*7	=0	R167
*1*6*8	=0	R168
*1*6*9	=0	R169
*1*7*0	=0	R170
*1*7*5	=0	R175
*1*7*7	=0	R177
*1*7*8	=0	R178
*1*7*9	=0	R179
*1*8*0	=0	R180
*1*8*1	=0	R181
*1*8*2	=0	APC2
*1*8*3	=0	APC2
*1*8*4	=0	APC1
*1*8*6	=0	APC2
*1*8*7	=0	APC1
*1*9*0	=0	APC2
*1*9*1	=0	APC1
*1*9*2	=0	APC2
*1*9*4	=0	APC2
*1*9*5	=0	APC1
*2*0*0	=0	APC2
*2*0*1	=0	APC1
*2*0*2	=0	R202
*2*0*5	=0	APC1
*2*0*6	=0	APC2
*2*0*7	=0	R207
*2*0*8	=0	R208
*2*0*9	=0	R209
*2*1*0	=0	R210
*2*1*1	=0	R211
*2*1*2	=0	R212
*2*1*3	=0	R213
*2*1*4	=0	R214
*2*1*5	=0	R215
*2*1*6	=0	R216
*2*2*0	=0	R220
*2*2*1	=0	R221
*2*2*2	=0	R222
*2*3*5	=0	R235
*2*3*6	=0	R236
*2*3*7	=0	R237

*2*3*8	=0	R238
*2*3*9	=0	R239
*2*4*0	=0	R240
*2*4*2	=0	R242
*2*4*4	=0	R244
*2*4*6	=0	R246
*2*4*8	=0	R248
*2*5*2	=0	R252
*2*5*7	=0	R257
*2*9*4	=0	R294
*2*9*5	=0	R295
*2*9*7	=0	R297
*2*9*8	=0	R298
*2*9*9	=0	R299
*3*0*0	=0	R300
*3*0*1	=0	R301
*3*0*3	=0	R303
*3*0*4	=0	R304
*3*0*5	=0	R305
*3*0*6	=0	R306
*3*0*7	=0	R307
*3*0*9	=0	R309
*3*1*0	=0	R310
*3*1*1	=0	R311
*3*1*4	=0	ASC
*3*1*5	=0	ASC
*3*1*6	=0	ASC
*3*1*7	=0	ASC
*3*1*8	=0	ASC
*3*1*9	=0	ASC
*3*2*0	=0	ASC
*3*2*1	=0	ASC
*3*2*2	=0	ASC
*3*2*3	=0	ASC
*3*2*4	=0	R324
*3*2*5	=0	R325
*3*2*6	=0	R326
*3*2*7	=0	R327
*3*2*8	=0	R328
*3*2*9	=0	R329
*3*3*0	=0	R330
*3*3*2	=0	R332
*3*3*3	=0	R333
*3*3*4	=0	R334
*3*3*9	=0	R339
*3*4*2	=0	R342
*3*4*4	=0	R344
*3*4*7	=0	R347
*3*5*0	=0	R350
*3*5*1	=0	R351
*3*5*2	=0	R352
*3*5*6	=0	R356
*3*6*3	=0	R363
*3*6*4	=0	R364
*3*6*9	=0	R369

	*3*7*1	=0	R371
	*3*7*2	=0	R372
	*3*7*4	=0	R374
	*3*7*5	=0	R375
	*3*7*7	=0	R377
	*3*8*6	=0	R386
	*3*8*7	=0	R387
	*3*8*8	=0	R388
	*3*8*9	=0	R389
	*3*9*0	=0	R390
	*3*9*1	=0	R391
	*3*9*2	=0	R392
	*3*9*3	=0	R393
	*3*9*4	=0	R394
	*3*9*5	=0	R395
	*3*9*7	=0	R397
	*3*9*9	=0	R399
	*4*0*1	=0	R401
	*4*0*3	=0	R403
	*4*0*6	=0	R406
	*4*0*7	=0	R407
	*4*1*7	=0	WNNT
	*4*1*8	=0	WNO
	*4*1*9	=0	WNNT
	*4*2*0	=0	WNPNT
	*4*2*1	=0	WNPO
	*4*2*2	=0	WNPNT
	*4*2*3	=0	WONT
	*4*2*4	=0	WOC
	*4*2*5	=0	WOPN
	*4*2*6	=0	WOP0
	*4*2*7	=0	R427
	*4*2*8	=0	R428
	*B	=0	STOP
*	\$1=0		EXIT
R2	TOP=CL*1		EXIT
R4	TOP=CL*3		EXIT
R6	TOP=CL*5		EXIT
R8	TOP=CL*7		EXIT
R10	TOP=CL*1*1		EXIT
R12	TOP=CL*1*3		EXIT
R14	TOP=CL*1*5		EXIT
R16	TOP=CL*1*7		EXIT
R25	AAP*1=WN+COM*1		EXIT
R26	AAP*1=COM*1		EXIT
R27	AAP*1=WOP+J		EXIT
R28	AAS=WN+COM*1		EXIT
R29	AAS=COM*1		EXIT
R30	AAS=WO+J		EXIT
R33	APN=PNBET		EXIT
R39	CL*1=PRE*0+*(+SUB*1X+*)		EXIT
R53	CL*3=PRE*1+*(+SUB*6X+*)		EXIT
R59	CL*5=PRE*3+*(+SUB*4X+*)		EXIT
R61	CL*7=PRE*4+*(+SUB*7W+*)		EXIT
R41	CL*1*1=PRE*1*1+*(+SUB*1*0+*)		EXIT

R43	CL*1*3=PRE*1*3+*(+SUB*1*1+*)	EXIT
R45	CL*1*5=PRE*1*5+*(+SUB*1*3+*)	EXIT
R47	CL*1*7=PRE*1*7+*(+SUB*1*6+*)	EXIT
R63	COL=SMC	EXIT
R64	COL=SMZ	EXIT
R65	COM*1=LP	EXIT
R66	COM*1=RF	EXIT
R67	COM*2=COM*1	EXIT
R68	COM*2=PNP*2	EXIT
R69	COM*3=COM*1	EXIT
R70	COM*3=SNP*2	EXIT
UNAT	Q=UN	QUADS
SQUAT	Q=S	*
QUADS	QUAD=CNJ	EXIT
TQUAT	Q=T	QUADS
FQUAT	Q=FQ	QUADS
UQUAT	Q=UQ	*
*	QUAD=PLY	EXIT
NQUAT	*(+Q=*--+*(+FQ	QUADS
R111	DUM*1+*(+\$,+,+NPH+*)=R*1+2+3+4+5+4+5+6	EXIT
R112	DUM*2=SNASN	EXIT
R113	DUM*2=PNAPN	EXIT
R114	DUM*2=PNBET	EXIT
R117	ES=SNP*2	EXIT
R119	F*1=G*1	EXIT
R120	F*2=G*2	EXIT
R121	MT+G*1=MORT	EXIT
*	G*1=RT	EXIT
R122	MT+G*1=MOLF	EXIT
*	G*1=LF	EXIT
R123	MT+G*1=MTOP	EXIT
*	G*1=TP	EXIT
R124	MT+G*1=MBOT	EXIT
*	G*1=BOT	EXIT
R125	MT+G*2=MCFN	EXIT
*	G*2=CFN	EXIT
R126	MT+G*2=MMID	EXIT
*	G*2=MID	EXIT
R130	IA=J*1	EXIT
R131	IA=J*2	EXIT
R132	*B+IA+*(+\$,+)*+C=1+1+J*1+3+4+5+6+CNJ+1+J*2+3+4+5+6+6	EXIT
R134	J=J*1	EXIT
R135	J=J*2	EXIT
R136	J*1=BG	EXIT
R137	J*1=LT	EXIT
R138	J*1=LG	EXIT
R139	J*1=SM	EXIT
R140	J*2=BK	EXIT
R141	J*2=WH	EXIT
R142	JER=BGR	EXIT
R143	JER=LTR	EXIT
R144	JER=LGR	EXIT
R145	JFR=SMR	EXIT
R146	LL=Z*1	EXIT
R147	LL=Z*2	EXIT

R148	LP=LL		EXIT
R152	N=TR		EXIT
R153	N=SQ		EXIT
R154	N=CIR		EXIT
R156	NO=NPLUR		EXIT
R158	NPH=SNP*2		EXIT
R159	NPH=PNP*2		EXIT
R160	NPLUR=TR		EXIT
R161	NPLUR=SQ		EXIT
R162	NPLUR=CIR		EXIT
R164	*R+NT=*-+1		EXIT
*	NT=*-		EXIT
P166	P=IN		EXIT
R167	P=NR		EXIT
R168	P=BEL		EXIT
R169	P=AB		EXIT
R170	P=TCH		EXIT
R175	P*3=BET		EXIT
R177	PA*1+*(+\$+*)=*B+IA+*(+3+*)+*C+CNJ+*B+AAP*1+2+3+4+*C		EXIT
R178	PA*1=AAP*1		EXIT
R179	PA*1=IA		EXIT
R180	PA*2+*(+\$+*)=*B+IA+*(+3+*)+*C+CNJ+*B+AAP*1+2+3+4+*C		EXIT
R181	PA*2=IA		EXIT
APC1	PC=NV		APC1
*	\$	//BACH A,OUT A	FUNC
APC2	PC=NV		APC2
*	\$	//BACH B,OUT A	FUNC
BACH	A		BACHA
	B		BACHB
BACHA	Q+\$+QUAD=EQ+2+CNJ		BACHB
RACHB	NV=	//*X3	CBACH
EQUA	X+*(+X+*)=2+1+*=+3+4		EQUA
*	\$	//*X3	*
*	\$1=1/.11	//*X3	EXIT
CBACH	\$1=1+1	//*X3	*
*	NV=	//*N3 1	BACHB
R202	PC*4S*6=PC		EXIT
R207	PC*6S=PC		EXIT
R208	\$1+*(+PNAPN+*)=1+2+PNBET+,+APN+4		EXIT
R209	PNBET=NV		R209
*	\$	//OUT B	FUNC
XOB	NV=PC		XOB
*	\$		EXIT
R210	PNBET=PC		EXIT
R211	PNP*1*2=NV		R211
*	\$	//OUT C	FUNC
XOC	NV=V		XOC
*	\$		EXIT
R212	PNP*1*2=PC		EXIT
R213	PNP*1*2=NV		R213
*	\$	//OUT B	FUNC
R214	PNP*2=NV		R214
*	\$	//OUT B	FUNC
R215	PNP*2=PC		EXIT
R216	PNP*2=NV		R216

*	\$	//OUT C	FUNC
R220	PNP*2A=NV		R220
*	\$	//OUT B	FUNC
R221	PNP*2A=PC		EXIT
R222	PNP*2A=NV		R222
*	\$	//OUT C	FUNC
R235	PNP*6S=NV		R235
*	\$	//OUT C	FUNC
R236	PNP*6S=PC*4S*6		EXIT
R237	PNP*6S=NV		R237
*		//OUT B	FUNC
R238	PPLB+*(+\$+*)=P*3+2+3+,+DUM*2+4		EXIT
R239	PRE*0=VP*0		EXIT
R240	PRF*1=VP*1		EXIT
R242	PRF*1*1=VP*1*1		EXIT
R244	PRF*1*3=VP*1*3		EXIT
R246	PRE*1*5=VP*1*5		EXIT
R248	PRF*1*7=VP*1*7		EXIT
R252	PRE*3=VP*3		EXIT
R257	PRE*4=VP*4		EXIT
R294	R=R*1		EXIT
R295	R=R*2		EXIT
R297	R=R*4		EXIT
R298	R=R*5		EXIT
R299	R=R*6		EXIT
R300	R*1=I L		EXIT
R301	R*2=JFP		EXIT
R303	R*4=P		EXIT
R304	R*5+*(+\$+*)=MT+DUM*1+2+3+4		EXIT
R305	R*6=COL		EXIT
R306	RE+*(+\$+*)=R+2+3+,+NPH+4		EXIT
R307	RF=PPLB		EXIT
R309	SA+*(+\$+*)=*B+IA+2+3+4+*C+CNJ+*B+AAS+2+3+4+*C		EXIT
R310	SA=IA		EXIT
R311	SA=AAS		EXIT
ASC	SC=NV		ASC
*	\$	//OUT D ,BACH R	FUNC
R324	\$1+*(+SNASN+*)=1+2+SNP*2+,+FS+4		EXIT
R325	SNP*1=NV		R325
*	\$	//OUT F	FUNC
XOF	NV=SC		XOF
*	\$		EXIT
R326	SNP*1=SC		EXIT
R327	SNP*1*3=SC		EXIT
R328	SNP*1*3=NV		R328
*	\$	//OUT F	FUNC
R329	SNP*1*6=SC		EXIT
R330	SNP*1*6=NV		R330
*	\$	//OUT F	FUNC
R332	SNP*2=SC		EXIT
R333	SNP*7=SC		EXIT
R334	SNP*7=NV		R334
*	\$	//OUT F	FUNC
R339	SUB*1=PNP*2A		EXIT
R342	SUB*1*0=PNP*2		EXIT

R344	SUB*1*1=PNP*1*2	EXIT
R347	SUB*1*3=SNP*1*3	EXIT
R350	SUB*1*6=SNP*1*6	EXIT
R351	SUB*1X=SUB*1Y	EXIT
R352	SUB*1Y=SUB*1	EXIT
R356	SUB*3=SNP*1	EXIT
R363	SUB*4X=SUB*4Y	EXIT
R364	SUB*4Y=SUB*3	EXIT
R369	SUB*6=PNP*6S	EXIT
R371	SUB*6X=SUB*6Y	EXIT
R372	SUB*6Y=SUB*6	EXIT
R374	SUB*7=SNP*7	EXIT
R375	SUB*7W=SUB*7Y	EXIT
R377	SUB*7Y=SUB*7	EXIT
R386	V=NV	R386
*	\$	FUNC
XOF	NV=V*1	XOF
*	\$	EXIT
R387	V*1=V*2	EXIT
R388	V*1=V*3	EXIT
R389	V*2=V*4	EXIT
R390	V*3=NV	R390
*	\$	FUNC
XOG	NV=	ZOG
*	\$=	*
*	\$1=1/.I1	EXIT
ZOG	\$=1+1	*
*	NV=	XOG
R391	V*4=NV	R391
*	\$	FUNC
R392	V*4=NV	R392
*	\$	FUNC
R393	V*5=V*3	EXIT
R394	*B+VP*0=*-+1+COM*1	EXIT
*	VP*0=*-+COM*1	EXIT
R395	VP*1=COM*1	EXIT
R397	*B+VP*1*1=*-+1+COM*2	EXIT
*	VP*1*1=*-+COM*2	EXIT
R399	VP*1*3=COM*2	EXIT
R401	*B+VP*1*5=*-+1+COM*1	EXIT
*	VP*1*5=*-+COM*3	EXIT
R403	VP*1*7=COM*3	EXIT
R406	*B+VP*3=*-+1+COM*3	EXIT
*	VP*3=*-+COM*3	EXIT
R407	VP*4=COM*1	EXIT
WNNT	WN=NT	EXIT
WNO	WN=0	EXIT
WNPNT	WNP=NT	EXIT
WNPO	WNP=0	EXIT
WONT	WO=NT	EXIT
WOO	WO=0	EXIT
WOPN	WOP=NT	EXIT
WOP0	WOP=0	EXIT
R427	Z*1=F*1	EXIT
R428	Z*2=F*2	EXIT

ZZ2	\$	//*Q2 1,*A6 1	ZZ1
FUNC	\$+NV	//*Q2 1	*
ZY1	\$+NV	//*Q5 1 2	ZY1
*	\$	//*X5	*
ZY2	*B+\$+*C=*D+2+*E		ZY2
*	\$	//*X2	*
ZZ1	\$+*B+\$+*B	//*Q6 1 2 3	ZZ1
*	\$+*B	//*Q6 1	*
*	*B+\$+*C=*7+*D+2+*E	//*A6 1	771
*	\$=1+*A	//*A2 2	*
*	*B+\$+*C=*D+2+*F		ZZ2
*	\$	//*X5	*
ZZ5	\$+*C	//*Q2 1 2,*X2	*
*	*F+\$+*C=*D+2+*E+*A	//*A2 4	Z75
*	\$=*Z+1	//*A5 1	*
*	\$1+\$+*B+\$+*C=1+2+*D+4+*E	//*Q5 1 2 3 4 5,*A2 1	ZZ5
*	\$=*A+*M+1+*N+*A	//*A6 1,*A2 5	*
DOF	*D+\$+*E=*B+2+*C		DOF
OUT A	*M+\$+*N=*B+*(+Q+NV+*))+*B+*B+NPLUR+*(+NV+*))+*C+QUAD+2+-		
	*C+*C		BACH
R	=*B+*B+PA*1+*(+PC+*))+*C+CNJ+2+*C		XOB
C	=*B+*B+PA*2+*(+V+*))+*C+CNJ+2+*C		XOC
D	=*B+*(+Q+NV+*))+*B+*B+N	+*(+NV+*))+*C+QUAD+2+-	
	*C+*C		BACH
F	=*B+*B+SA+*(+SC+*))+*C+CNJ+2+*C		XOF
F	=*B+2+CNJ+*B+NO+*(+V*1+*))+*C+*C		XOF
G	=*B+2+CNJ+*B+NPLUR+*(+NV+*))+*C+*C		XOG
H	=*B+2+CNJ+*B+NPLUR+*(+NV+*))+*C+1+2+3+*C		XOH
I	=*B+2+CNJ+*B+NPLUR+*(+NV+*))+*C+1+2+3+*C		XOI
XOH	*M+\$+NV+\$+*N=1+2+V*5+4+5		XOH
*	*M+\$+*N=2		XOG
XOI	*M+\$+NV+\$+*N=1+2+V*2+4+5		XOI
*	*M+\$+*N=2		XOG
STOP	*B+\$+*C=*(+2+*))		STOP
*	\$	//*WSL 1	*
*	\$=*A+*A+*A+*A+*A+*A	//*A1 1,-	
*A2 2,*A3 3,*A4 4,*A5 5,*A6 6			*
*	\$=0		STAR
FINISH			*
END			

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