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The Numerical Solution of A Nonseparable Elliptic Partial Differential Equation By Preconditioned Conjugate Gradients

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U. S. DEPARTMENT OF COMMERCE
National Bureau of Standards
National Engineering Laboratory
Center for Applied Mathematics
Washington, DC 20234

April 1980

Final Report

Prepared for:

for Applied Mathematics
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**THE NUMERICAL SOLUTION OF A
NONSEPARABLE ELLIPTIC PARTIAL
DIFFERENTIAL EQUATION BY
PRECONDITIONED CONJUGATE
GRADIENTS**

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U.S. DEPARTMENT OF COMMERCE, Malcolm Baldrige, *Secretary*
NATIONAL BUREAU OF STANDARDS, Ernest Ambler, *Director*

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The Numerical Solution of a Nonseparable Elliptic Partial Differential Equation by Preconditioned Conjugate Gradients

John Gregg Lewis* and Ronald G. Rehm†

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In this report the combination of an iterative technique, the conjugate gradient algorithm, with a fast direct method, cyclic reduction, is used to solve the linear algebraic equations resulting from discretization of a nonseparable elliptic partial differential equation. An expository discussion of the conjugate gradient and preconditioned conjugate gradient algorithms and of their use in the solution of partial differential equations is presented. New results extending the use of the preconditioned conjugate gradients technique to singular linear equations which arise from discretized elliptic equations with Neumann boundary conditions are also given. The algorithms are applied to solve a specific elliptic equation which arises in the study of buoyant convection produced by a room fire. A code was developed to implement the algorithms for this application. Numerical results obtained through testing and use of the code are discussed.

Key Words: Conjugate gradient algorithm; elliptic partial differential equations; iterative methods for linear algebraic equations; Neumann boundary conditions; sparse matrices.

The numerical solution of the linear algebraic equations which result from a discretization of an elliptic partial differential equation has been, and continues to be, the focus of much research in numerical analysis. Over the past 25 years there have been many advances, some toward improved iterative schemes (ADI, SOR, etc.), others toward fast direct methods (most notably those based on FFT's). See Rice [1]¹ for a brief survey of the impact of these advances. In this paper we discuss the combination of an iterative technique, the conjugate gradient algorithm, with a fast direct method, cyclic reduction, to extend the capabilities of the fast solver. For examples of the use of similar combinations of algorithms, see Concus & Golub [2], Concus, Golub and O'Leary [3], O'Leary [4] and O'Leary and Widlund [5].

The work described in this paper resulted from a study of buoyant convection carried out at the National Bureau of Standards, [6, 7, 8]. The specific elliptic partial differential equation for pressure arising in this work is used as a model problem in the present paper. The experience of the first author with the buoyant convection model motivated the writing of sections 1 and 2, which gives an exposition of the conjugate gradient and preconditioned conjugate gradient algorithms. These sections contain no new material; they were written so that this paper may be accessible to an audience unfamiliar with the development of the conjugate gradient algorithm and its use in the solution of partial differential equations. Section 3 has a discussion of the model problem, the pressure equation. Section 4 contains several new results extending the use of preconditioned conjugate gradient technology to the singular linear equations which result from Neumann boundary value problems. We conclude with some numerical examples from the buoyant convection problem. A listing of a FORTRAN program implementing the algorithm discussed here is presented in [15].

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¹ Figures in brackets indicate literature references at the end of this paper.

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1. An introduction to the conjugate gradient algorithm

The discretization of an elliptic differential operator by standard finite difference or finite element techniques produces, as the finite analog of the continuous operator, very large matrices with most elements zero. From an algebraic standpoint we can very often view the numerical solution of the differential equation as:

$$\text{Solve } Ax = b \quad (1)$$

where A is a k by k , (k large) real symmetric, positive definite and sparse matrix.

Fast direct methods with minimal storage requirements exist when A is the finite difference operator resulting from a separable elliptic operator on a rectangular region, and for some other common, but specific, cases. If the operator is nonseparable, or the region non-rectangular, the special direct methods may not apply. Moreover, the use of general direct methods such as factoring A into a Cholesky decomposition,

$$A = LL^T,$$

often impose unbearable storage requirements because many zero entries are filled in during the decomposition.

Iterative methods, such as SOR and Gauss-Seidel, solve a transformation of the problem. Instead of solving (1) directly, they "split" A as $A = M - N$, and solve the equivalent problem through an iteration of the form $Mx^{k+1} = Nx^k + b$; i.e., find a fixed point of the equation

$$Mx = Nx + b. \quad (2)$$

The efficiency of the solution method depends in part on the appropriateness of the splitting and in part on acceleration of the iteration toward the fixed point.

The conjugate gradient algorithm can be motivated as the solution of another transformation of problem (1). Consider the inner product $\langle x, y \rangle \equiv x^T A y$. This induces a norm $\|x\|_A \equiv \sqrt{x^T A x}$ on R^k when A is positive definite.

Problem (1) is equivalent to: find x to minimize

$$E(x) = (x - x^*)^T A(x - x^*) = \|x - x^*\|_A^2 \quad (3)$$

where x^* is the solution of (1). The problems are equivalent because $E(x) \geq 0$ for all x , and $E(x) = 0 \Leftrightarrow x = x^*$. The immediate usefulness of the transformation is not obvious, especially since computing $E(x)$ appears to require the solution, x^* , of the problem we wish to solve. Note however, that we can evaluate whether a given x is or is not a solution by computing the residual $r = b - Ax$. Further, while we cannot evaluate $E(x)$ directly, we can evaluate its gradient vector:

$$\nabla E(x) = 2(Ax - Ax^*) = 2(Ax - b) = -2r.$$

Hence, r gives us the direction in which E decreases most rapidly (unless r is identically zero, which characterizes the solution).

These two characteristics of problem (3) lead directly to a simple algorithm for solving (3):

STEEPEST DESCENT: Given a guess x_i , not a solution,

$$\left\{ \begin{array}{l} \text{set } r_i = b - Ax_i \\ \text{set } x_{i+1} = x_i + \alpha_i r_i \end{array} \right.$$

where α_i is a scalar chosen to make x_{i+1} minimize $E(x_{i+1})$ for all vectors of the form $x_i + \beta r_i$.

We can easily compute the gradient r_i . It is also easy to compute α_i so that $E(x_{i+1}) = E(x_i + \alpha_i r_i) \leq E(x_i + \beta r_i)$ for all β . The minimum of $E(x_i + \beta r_i)$ is given by the solution of

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$$\frac{d}{d\beta} E(x_i + \beta r_i) = 0$$

But

$$\frac{d}{d\beta} E(x_i + \beta r_i) = 2(\beta r_i^T A r_i - r_i^T r_i),$$

so

$$\alpha_i = \frac{r_i^T r_i}{r_i^T A r_i} = \frac{\|r_i\|_2^2}{\|r_i\|_A^2} \quad (4)$$

is the optimal choice for α_i . Hence, we can carry out the iteration to minimize $E(x)$ without computing E . Note also that A enters the iteration only in forming the products Ax_i and $A r_i$.

Unfortunately steepest descent is not a practical algorithm because the convergence can be very slow. We can obtain more rapid convergence by using not the steepest descent directions $\{r_i\}$, but instead a sequence of downhill or descent directions $\{p_i\}$ which satisfy additional properties. We characterize the term "descent direction" and simultaneously find the optimal distance to move in that direction, through the following simple result.

LEMMA: If $p^T r \neq 0$, then

$$\beta = \frac{p^T r}{p^T A p} \text{ minimizes } E(x + \beta p) \text{ for all } \beta. \quad (5)$$

PROOF: Simply differentiate $E(x + \beta p)$ with respect to β :

$$\begin{aligned} \frac{d}{d\beta} E(x + \beta p) &= 2(p^T A(x - x^*) + \beta p^T A p) \\ &= 2(-p^T r + \beta p^T A p). \end{aligned}$$

It is a simple computation to show that $E(x + \beta p) < E(x)$. Note that β is positive if $p^T r > 0$, which characterizes p as a descent, not an ascent, direction.

The first observation which leads to improved convergence is that the choice in (5) for β , not only solves the one-dimensional minimization, it gives us the p -component of x^* exactly. To see this, note that A positive definite implies that we can uniquely write

$$\begin{aligned} x + \beta p &= \alpha p + \delta w \\ x^* &= \alpha^* p + \delta^* z \end{aligned} \quad (6)$$

where $w^T A p = 0 = z^T A p$. This last condition is read: w, z are A -orthogonal or A -conjugate to p . Now substitute the decompositions (6) into E:

$$E(x + \beta p) = (\alpha - \alpha^*)^2 p^T A p + \text{terms not involving } \alpha, \alpha^* \text{ or } p \quad (7)$$

The choice (5) for β minimizes the left side of (7), and clearly $\alpha = \alpha^*$ minimizes the right hand side of (7). Hence choosing $\beta = \frac{p^T r}{p^T A p}$ implies that

$$(x + \beta p) - x^* = \delta w - \delta^* z. \quad (8)$$

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The error vector is A -conjugate to p , i.e., lies in the $k - 1$ dimensional subspace of vectors t such that $t^T A p = 0$. If all succeeding descent directions are chosen from this subspace, we can preserve the exact solution of the problem in the direction p .

The second observation is that we can easily choose successive directions p_1, p_2, p_3, \dots such that the $\{p_i\}$ are all pairwise A -conjugate descent directions (hence are "A-conjugate gradients") and hence, such that the successive errors $x_1 - x^*, x_2 - x^*, x_3 - x^*$ are constrained to successively smaller dimensional subspaces. The set $\{p_i\}, i = 1, \dots, k$ gives a basis for R^k and $x_k - x^*$ must be A -conjugate to all of the $\{p_i\}$, hence must be zero. Thus, the process must give the exact answer after at most k iterations. We now write out the conjugate gradient iteration directly:

CONJUGATE GRADIENTS: Given x_1 , compute $r_1 = b - Ax_1$. Set the initial descent direction p_1 to be r_1 . For $i = 2, 3, 4, \dots$, do.

Move to the minimum in direction p_i and evaluate the new residual

$$\begin{aligned} \alpha_i &= \frac{r_i^T p_i}{p_i^T A p_i} \\ x_{i+1} &= x_i + \alpha_i p_i \\ r_{i+1} &= b - Ax_{i+1} \\ &= r_i - \alpha_i A p_i \end{aligned} \tag{9}$$

Compute a new descent direction A -orthogonal to all of its predecessors

$$\begin{aligned} \beta_i &= -\frac{r_{i+1}^T A p_i}{p_i^T A p_i} \\ p_{i+1} &= r_{i+1} + \beta_i p_i \end{aligned} \tag{10}$$

The initial step (9) in the iteration is the step in the steepest descent direction p_i , as discussed above. Step (10), the choice of a new descent direction, is simply a single step of Gram-Schmidt orthogonalization, to remove from the steepest descent direction its component in the direction $A p_i$.

Hence

$$p_{i+1}^T A p_i = 0 \tag{11}$$

Further the choice of α_i implies directly that

$$r_{i+1}^T p_i = 0 \tag{12}$$

The following identities are derived easily from (9) and (10):

$$\begin{aligned} r_{i+1}^T r_i &= 0; \\ r_i^T p_i &= r_i^T r_i; \\ r_i^T A p_i &= p_i^T A p_i. \end{aligned} \tag{13}$$

The most important algebraic consequences of (9) and (10) yield only to a lengthy inductive argument (not given here—see Reid [12] or Hestenes and Stiefel [13] for example)

$$\begin{aligned} r_{i+1}^T r_j &= 0 \\ p_{i+1}^T A p_j &= 0 \end{aligned} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} \text{for all } j \leq i \tag{14}$$

The identities in (14) imply that all of the $\{p_i\}$ are A -conjugate, even though we perform an explicit orthogonalization only to one descent direction in (10).

In the conjugate gradient iteration observe that A enters only in forming a matrix-vector product, Ap_i ; there are no transformations done which could destroy the sparseness of A . The dominant cost per step is usually that of forming the product Ap_i ; in this case the conjugate gradient iteration is very little more expensive per iteration than the steepest descent algorithm and it offers the guarantee of convergence within k steps.

The practical reason for using the conjugate gradient algorithm is that we often obtain satisfactory accuracy after only a very few iterations. Certain theoretical bounds for the convergence of the algorithm depend on the condition number, κ , of A , defined by:

$$\kappa = \frac{\lambda_{max}(A)}{\lambda_{min}(A)} \quad (15)$$

where $\lambda_{max}(A)$ is the largest eigenvalue of A and $\lambda_{min}(A)$ is the smallest eigenvalue of A (Note: A is positive definite, so κ is positive, and at least as great as one). The following bounds can be found in the literature (see Daniel [14], for example):

$$\|x_i - x^*\|^2 \leq 4 \frac{E(x_i)}{\lambda_{min}} \cdot \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^{2(i-1)}$$

$$E(x_i) \leq 4 \left(\frac{1 - \sqrt{\frac{1}{\kappa}}}{1 + \sqrt{\frac{1}{\kappa}}} \right)^{2(i-1)} E(x_1)$$

Clearly, convergence is rapid when κ is close to one (and takes place in one step if κ equals one). For well-conditioned problems very few iterations will suffice to obtain highly accurate solutions.

2. Preconditioned Conjugate Gradients and Matrix-Splittings.

Our basic problem is to solve

$$Ax = b, \quad (16)$$

which is the large, sparse linear system which arises from the discretization of an elliptic partial differential equation. We can assume that the matrix A is symmetric and positive definite, so the method of conjugate gradients certainly can be used. However, the condition number, $\kappa(A)$, is usually very large for these problems; the conjugate gradient algorithm converges very slowly when applied directly to A and is not competitive with other methods.

The preconditioned conjugate gradients method arises, like SOR, from matrix splittings. Consider

$$A = M - N$$

where M is another positive-definite, symmetric matrix, but one for which we can easily solve linear systems. There exists a symmetric positive-definite matrix $M^{-1/2}$ such that

$$M^{-1} = (M^{-1/2})(M^{-1/2})$$

or

$$[(M^{-1/2})(M^{-1/2})]^{-1} = M$$

(This is the symmetric matrix whose eigenvectors are the same as M 's and whose corresponding eigenvalues are the reciprocals of the square roots of the eigenvalues of M . We shall not use this form, however; we need only the formal existence of $M^{-1/2}$.) We can solve the equations $Ax = b$ by computing

$$d = M^{-1/2}b; \quad (17)$$

and solving

$$Cz = d \quad (18)$$

where $C = M^{-1/2}AM^{-1/2}$, by the conjugate gradient algorithm (Note that C is a positive-definite, symmetric matrix), finally computing

$$x = M^{-1/2}z \quad (19)$$

This transformation from $Ax = b$ to $Cy = d$ will be an effective method if:

- a) the condition number, $\kappa(C)$, is close to 1 so that the conjugate gradient algorithm converges rapidly, and
- b) the multiplication $C \cdot y$ for any vector y can be computed efficiently and sparsely.

We shall examine condition (a) first. Note that

$$A = M - N$$

implies

$$C = M^{-1/2}AM^{-1/2} = I - (M^{-1/2}NM^{-1/2}) = I - R.$$

Hence

$$\kappa(C) = \frac{\lambda_{\max}(I-R)}{\lambda_{\min}(I-R)} \geq 1$$

This condition number, $\kappa(C)$, will be close to 1 if $\lambda_{\max}(R)$ is small, which is true if all of the elements of R are small. (The condition number is 1 exactly when R is 0, or when $M = A$). We want, then, to choose M so that M represents A as well as possible (makes R as small as possible), and such that the choice for M allows for condition (b). We discuss a specific choice for M in section 3, the discussion of the model problem. Other examples are given in Concus, Golub, and O'Leary [3], O'Leary [4], and Meijerink and van der Vorst [9].

We now turn to condition (b); with a formal derivation of the actual algorithm used as "Preconditioned conjugate gradients," we show that condition (b) is met whenever M is such that the linear equation

$$Mz = w$$

can be solved efficiently and sparsely. Reconsider the conjugate gradient iteration to solve $Cz = d$, given an initial guess z_1 :

$$(CG_1): \text{ set } \bar{r}_1 = d - Cz_1;$$

$$\text{ set } \bar{p}_1 = \bar{r}_1;$$

$$\text{ for } i = 2, 3, 4, \dots$$

$$\alpha_i = \frac{\bar{r}_i^T \bar{r}_i}{\bar{p}_i^T C \bar{p}_i}$$

$$z_{i+1} = z_i + \alpha_i \bar{p}_i$$

$$\bar{r}_{i+1} = \bar{r}_i - \alpha_i C \bar{p}_i$$

$$\beta_i = \frac{\bar{r}_{i+1}^T \bar{r}_{i+1}}{\bar{r}_i^T \bar{r}_i}$$

$$\bar{p}_{i+1} = \bar{r}_{i+1} + \beta_i \bar{p}_i$$

In the above, the residual vectors $\{\bar{r}_i\}$ and descent directions $\{\bar{p}_i\}$ have bars on them to denote that they pertain to the scaled problem $Cz = d$. Now view (CG_c) from the original space in which we solve $Ax = b$. Define

$$x_i = M^{-1/2} z_i,$$

$$r_i = b - Ax_i$$

Then

$$\bar{r}_i = d - Cz_i = M^{-1/2} r_i$$

Similarly, we can define formally

$$p_i \equiv M^{-1/2} \bar{p}_i$$

Then

$$\alpha_i = \frac{\bar{r}_i^T \bar{r}_i}{\bar{p}_i^T C \bar{p}_i} = \frac{r_i^T (M^{-1/2})^T (M^{-1/2}) r_i}{p_i^T (M^{1/2})^T C M^{1/2} p_i} = \frac{r_i^T (M^{-1} r_i)}{p_i^T A p_i}$$

Updating z_i corresponds to taking

$$x_{i+1} = x_i + \alpha_i M^{-1/2} \bar{p}_i = x_i + \alpha_i p_i$$

The corrected residual is

$$\begin{aligned} r_{i+1} &= M^{1/2} \bar{r}_{i+1} \\ &= M^{1/2} (\bar{r}_i - \alpha_i C \bar{p}_i) \\ &= M^{1/2} \bar{r}_i - \alpha_i A M^{-1/2} \bar{p}_i \\ &= r_i - \alpha_i A p_i. \end{aligned}$$

The Gram-Schmidt coefficient β_i can be computed as:

$$\beta_i = \frac{\bar{r}_{i+1}^T \bar{r}_{i+1}}{\bar{r}_i^T \bar{r}_i} = \frac{r_{i+1}^T (M^{-1} r_{i+1})}{r_i^T (M^{-1} r_i)}$$

The new direction becomes:

$$p_{i+1} = M^{-1/2} \bar{p}_{i+1} = M^{-1/2} (\bar{r}_{i+1} + \beta_i \bar{p}_i)$$

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$$\begin{aligned}
&= M^{-1/2}(M^{-1/2} r_{i+1}) + \beta_i p_i \\
&= M^{-1} r_{i+1} + \beta_i p_i.
\end{aligned}$$

The result of these manipulations is that we can now write a conjugate-gradient like iteration to solve the equations $Ax = b$, but whose convergence rate depends on

$$C = M^{-1/2} A M^{-1/2}, \text{ not } A.$$

PRECONDITIONED CONJUGATE-GRADIENTS (PCG): Given initial guess x_1 : compute

$$v_1 = Ax_1$$

$$r_1 = b - Ax_1 = b - v_1$$

$$\text{solve } Mu_1 = r_1$$

$$\text{and set } p_1 = u_1$$

For $i = 2, 3, 4, \dots$

$$\text{compute } v_i = Ap_i;$$

$$\text{set } \alpha_i = \frac{r_i^T u_i}{p_i^T v_i} = \frac{r_i^T M^{-1} r_i}{p_i^T A p_i}$$

$$\text{update } x_{i+1} = x_i + \alpha_i p_i$$

$$r_{i+1} = r_i - \alpha_i v_i (= r_i - \alpha_i A p_i)$$

$$\text{solve } Mu_{i+1} = r_{i+1}$$

$$\text{set } \beta_i = \frac{r_{i+1}^T u_{i+1}}{r_i^T u_i} (= \frac{r_{i+1}^T M^{-1} r_{i+1}}{r_i^T M^{-1} r_i})$$

$$\text{finally, } p_{i+1} = u_{i+1} + \beta_i p_i$$

We introduced new vectors $\{u_i\}$ and $\{v_i\}$ in the algorithm above to emphasize the fact that the matrices A and M appear only one time during each iteration, and in very specific ways. The matrix A appears only in the formation of the product Ap_i (or Ax_i), an operation which can be done very efficiently for most representations of sparse matrices. The matrix M appears only implicitly; we must be able to solve the linear systems $Mu = r$ efficiently and in little storage, and this is the primary restriction on M . Note that the matrix square-roots, $M^{1/2}$, do not appear at all. The purpose of the matrix M is to *scale* or *pre-condition* the problem so that convergence takes place very quickly. The cost, of course, is that each iteration now requires the solution of a linear equation as well as the formation of a matrix-vector product.

3. A model problem—the pressure equation.

The problem to be discussed here arises in a model of buoyant convection (Rehm and Baum [6]). Specifically, at each time step in the solution of a mixed hyperbolic-elliptic system, we are required to solve:

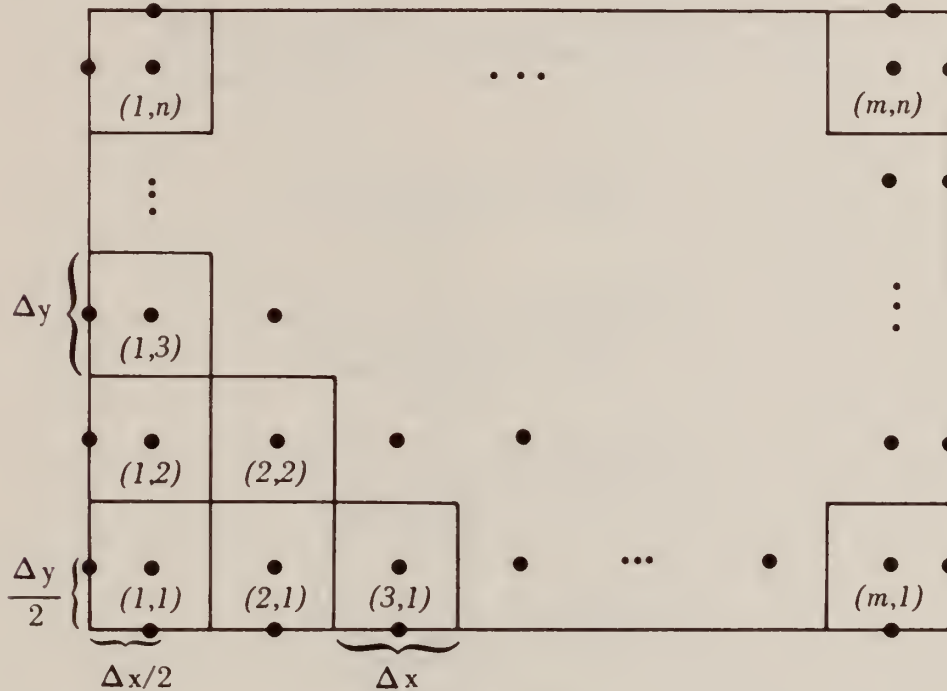
$$\nabla \cdot \left(\frac{1}{\varrho(x,y)} \nabla P(x,y) \right) = f(x,y) \quad (20)$$

on a rectangle R subject to the following condition on the exterior normal derivative

$$\frac{\partial P}{\partial \eta}(x,y) = \varrho(x,y) \cdot g(x,y) \text{ on the boundary } \partial R.$$

Here ρ denotes gas density and P the unknown pressure. We assume that ρ depends on both spatial variables. Equation (1) then describes a *nonseparable* elliptic equation. Since we must solve this equation repeatedly, speed is paramount. Were this a separable equation, the discrete linear system discussed below could be solved directly by the cyclic reduction routines of Sweet and Schwarztrauber [10]. Were the density and its first and second derivatives known analytically, the problem could be converted to an ordinary Poisson equation by the techniques of Concus and Golub [2]. Neither of these conditions can be met, which forces us to consider the discrete linear system in more detail.

In the context of the buoyant convection problem, we are given the density ρ , and g and f , only on a discrete set of points; the solution P will be produced on the same grid of discrete points. A grid suitable for the hydrodynamic calculations is:



The functions ρ and f are given on the interior points; g is given at the circles on the boundary. Note that the interior grid is offset by one half grid spacing from the boundary.

The derivatives in (20) are replaced by second-order accurate centered finite differences. For example,

$$\frac{\partial}{\partial x} \cdot \left(\frac{1}{\rho(x,y)} \frac{\partial}{\partial x} P(x,y) \right)^{1/2}$$

becomes

$$\frac{1}{\Delta x} \left(\frac{1}{\rho(x + \frac{\Delta x}{2}, y)} \left[\frac{1}{\Delta x} (P(x + \Delta x, y) - P(x, y)) \right] - \frac{1}{\rho(x - \frac{\Delta x}{2}, y)} \left[\frac{1}{\Delta x} (P(x, y) - P(x - \Delta x, y)) \right] \right)$$

Note that we require the reciprocal of the density at an intermediate point: an $O(\Delta x)$ approximation will suffice to give second order accuracy for the operator. For consistency with the hydrodynamic equations in the buoyant convection model, we use the reciprocal of the average density:

$$\frac{1}{\varrho(x + \frac{\Delta x}{2}, y)} \cong \frac{2}{[\varrho(x + \Delta x, y) + \varrho(x, y)]}$$

Define $\varrho_{ij} = \varrho((i-1/2)\Delta x, (j-1/2)\Delta y)$, and similarly for P_{ij} , f_{ij} , and g_{ij} . Then the general equation for an interior grid point i, j not adjacent to the boundary is

$$\begin{aligned} & \left\{ \left[\underbrace{\left(\frac{-2}{\varrho_{i-1,j} + \varrho_{i,j}} \right)}_{d_1} + \underbrace{\left(\frac{-2}{\varrho_{i,j} + \varrho_{i+1,j}} \right)}_{d_2} \right] \left(\frac{1}{\Delta x} \right)^2 + \right. \\ & \left. \left[\underbrace{\left(\frac{-2}{\varrho_{i,j-1} + \varrho_{i,j}} \right)}_{d_3} + \underbrace{\left(\frac{-2}{\varrho_{i,j} + \varrho_{i,j+1}} \right)}_{d_4} \right] \left(\frac{1}{\Delta y} \right)^2 \right\} P_{ij} \\ & + \underbrace{\left(\frac{2}{\varrho_{i,j} + \varrho_{i-1,j}} \right)}_{-d_1} \left(\frac{1}{\Delta x} \right)^2 P_{i-1,j} \\ & + \underbrace{\left(\frac{2}{\varrho_{i,j} + \varrho_{i,j+1}} \right)}_{-d_2} \left(\frac{1}{\Delta x} \right)^2 P_{i+1,j} \\ & + \underbrace{\left(\frac{2}{\varrho_{i,j-1} + \varrho_{i,j}} \right)}_{-d_3} \left(\frac{1}{\Delta y} \right)^2 P_{i,j-1} \\ & + \underbrace{\left(\frac{2}{\varrho_{i,j} + \varrho_{i,j+1}} \right)}_{-d_4} \left(\frac{1}{\Delta y} \right)^2 P_{i,j+1} = f_{ij} \end{aligned} \tag{21_{ij}}$$

The adjustment for points adjacent to the boundary uses the second order centered approximation to the boundary derivative: on boundary k (left = 1, right = 2, lower = 3, upper = 4), the terms d_k are evaluated using the first interior mesh point and an image point (one-half grid spacing outside the region). Formally the terms d_k are evaluated at $i = 1/2, j$ at the left boundary for example. The discretized form of the Neuman boundary conditions become

$$\left(\frac{1}{\Delta x} \right)^2 \left(-\frac{1}{2} d_1 p_{1,j} + \frac{1}{2} d_1 p_{0,j} \right) = g_{1,1/2,j} / \Delta x \equiv g_1(j) \tag{22_1}$$

$$\left(\frac{1}{\Delta x} \right)^2 \left(-\frac{1}{2} d_2 p_{m,j} + \frac{1}{2} d_2 p_{m+1,j} \right) = g_{m+1/2,j} / \Delta x \equiv g_2(j) \tag{22_2}$$

$$\left(\frac{1}{\Delta y} \right)^2 \left(-\frac{1}{2} d_3 p_{i,1} + \frac{1}{2} d_3 p_{i,0} \right) = g_{i,1/2} / \Delta y \equiv g_3(i) \tag{22_3}$$

$$\left(\frac{1}{\Delta y} \right)^2 \left(-\frac{1}{2} d_4 p_{i,n} + \frac{1}{2} d_4 p_{i,n+1} \right) = g_{i,n+1/2} / \Delta y \equiv g_4(i) \tag{22_4}$$

For a point adjacent to boundary k , subtract equation (22_k) from equation $(21_{i,j})$ which eliminates the d_k terms, and with them, all references to points i, j outside the grid. The right hand side is replaced by

$$f_{ij} - g_k$$

Corner points are treated by subtracting (22_k) and (22_l) from $(21_{i,j})$, where the corner is adjacent to both its k -th and l -th boundary.

We arrange the grid points in the order $\{(1,1), (2,1), \dots, (m,1), (1,2), (2,2), \dots, (m,2), \dots, (1,n), \dots, (m,n)\}$ and write the equation for each point in this order. The result is the matrix equation

where T_i is an $m \times m$ symmetric tridiagonal matrix, and O_i is an $m \times m$ diagonal matrix. We shall denote this matrix equation by

$$Ax = b \tag{23}$$

where b represents the right hand side, f , of the pressure equation, adjusted by the boundary data, g .

This is the linear algebraic system we need to solve. We make several observations:

1. A is symmetric
2. A is sparse. Only five diagonals contain non-zero elements.
3. The eigenvalues of A are all real and non-positive; A is a negative semi-definite matrix, generally of rank $mn - 1$. Hence $-A$ is a positive semi-definite matrix.

Although we could apply the method of conjugate gradients directly to $-A$, convergence would be very slow. O'Leary [4] and Meijerink and Van der Vorst [9] suggest several different approaches for creating scaling matrices M to use a preconditioned conjugate gradient scheme. One general approach is to consider that A originated from a differential equation and let M be the discrete operator from a separable differential equation which approximates the pressure equation.

A specific choice is suggested by Concus and Golub [2]: Let L be the discretization of the Poisson equation

$$\nabla^2 P = f$$

with Neumann boundary conditions (on the same staggered grid). Then

We assume that we have available a good subroutine for solving the equation

$$Lu = v,$$

e.g., subroutine BLKTRI from the NCAR package of subroutines for elliptic partial differential equations [10].

Let $D^{1/2}$ be the diagonal matrix with entries

$$d_{kk} = \sqrt{a_{kk}/l_{kk}},$$

the square root of the ratio of corresponding main diagonal entries of A and L . Then

$$M = D^{1/2} L D^{1/2}$$

is a symmetric negative semi-definite sparse matrix whose main diagonal is identical to that of A . Hence, in the splitting

$$A = M - N,$$

the matrix N has zero main diagonal and non-zero entries on at most four off-diagonals.

In the preconditioned conjugate gradient iteration we must be able to solve

$$Mz = r$$

This is done without excessive storage demands by:

- a) compute $w = D^{-1/2} r$ (D is a diagonal matrix)
- b) solve $Lu = w$ by cyclic reduction (BLKTRI)
- c) compute $z = D^{-1/2} u$.

For computational convenience, to avoid the repeated multiplications by the diagonal matrix $D^{1/2}$ or its inverse, we replace the original problem

$$Ax = b$$

by

$$\hat{A}\hat{x} = \hat{b}$$

where

$$\hat{A} = D^{-1/2} A D^{-1/2}$$

$$\hat{b} = D^{-1/2} b$$

solve $Ax = b$ by preconditioned conjugate gradients with L as the scaling matrix, and finally compute

$$x = D^{-1/2} \hat{x}$$

This is the approach actually used in the computations reported in section 5.

4. The Preconditioned Conjugate Gradient Algorithm for a Semi-definite Matrix.

In our discussion of the discrete operators A and L in the previous section we have ignored one characteristic of these operators which renders them unsuitable for direct use in a conjugate gradient or preconditioned conjugate gradient iteration. The conjugate gradient algorithm requires that A be a positive-definite matrix; otherwise the function E may have zeroes other than at the solution of the linear equation. The

algorithm becomes transparently a maximization algorithm for negative-definite matrices. (There is no need to change signs implicitly or explicitly to solve linear systems involving negative definite matrices, e.g., the discrete Laplacian with Dirichlet boundary conditions). However, our operators A and L are both negative semidefinite; they have negative eigenvalues, and also, one zero eigenvalue each.

In this section we will extend the theory of the conjugate gradient algorithm to allow for the special case of a semi-definite matrix with known nullspace. We shall then extend the preconditioned conjugate gradient algorithm to allow both A and M to be of this special type. The first part of the section will be elementary for readers well-versed in the conjugate gradient theory; the second part contains some new and unexpected results.

The problem is simple: A is a singular matrix. This is shown easily by noting that whenever a term d_k appears on the diagonal of A , the opposite term $-d_k$ appears on one of the off-diagonals. Hence, the sum of the coefficients in any row of A is exactly zero. But this is the same as saying

$$Ae = 0$$

where e is the vector $(1,1,1, \dots, 1)^T$. In general, $Ax = 0 \iff x = \alpha e$, where α is a real scalar.

The singularity of A is related to the Neumann boundary conditions for the pressure equation. If P is a solution to the differential equation so is $P + c$, where c is any constant. Similarly, if x is any solution to (3), so is $x + \alpha e$, where α is any scalar. But this is equivalent to adding the constant α to each point x_{ij} of the solution. The singularity of the operators also implies that not every system of equations has a solution. A system of equations is consistent if and only if it has a solution. The characterization of consistency for these operators is simple to express: The pressure equation is consistent \iff

$$\iint_R f = \int_{\partial R} g \quad (25)$$

The linear equations $Ax = b$ are consistent \iff

$$\begin{aligned} b^T e = 0, \text{ that is, when } b \text{ in Range}(A). \text{ But} \\ b^T e = 0 \iff \sum_{ij} b_{ij} = 0 \\ \iff \sum_{ij} f_{ij} = \sum_j (g_1(j) + g_2(j)) + \sum_i (g_3(i) + g_4(i)) \end{aligned}$$

the discrete analog of the integral equality (25)

Further, even when the equation is consistent, the solution is not unique. There is a unique solution of shortest length in the usual L_2 norm. For the continuous case it is the solution with mean pressure zero, i.e.,

$$\iint_R P = 0.$$

In the discrete case the analog is

$$x^T e = 0;$$

again, the mean (discrete) pressure is zero.

We shall now discuss the modifications to the conjugate gradient algorithm which would be necessary to obtain this unique solution to a consistent system of equations. For generality, assume that we want to solve

$$Ax = b \quad (26)$$

where A is symmetric positive semi-definite of dimension k with known nullspace the span of $\{n\}$. Hence, the rank of A is one less than its dimension. Further, assume $b^T n = 0$, so that (26) is a consistent system.

(note: if (26) is not consistent, $\hat{b} = b - \frac{b^T n}{n^T n} n$ satisfies the consistency condition, and any solution of $Ax = \hat{b}$ is a least squares solution of (26)).

Let the vectors $\{n, q_1, q_2, q_3, \dots, q_k\}$ form an orthonormal basis for R^k . Then $\text{Range}(A) = \text{span}\{q_1, q_2, \dots, q_k\}$

Let Q be the orthogonal matrix ($Q^T = Q^{-1}$)

$$Q = [q_1 \ q_2 \ q_3 \ \dots \ q_k]$$

Then

$$Ax = b \iff (Q^T A Q)(Q^T x) = Q^T b \quad (27)$$

But, by the consistency condition $b^T n = 0$

$$Q^T b = \begin{bmatrix} 0 \\ d_1 \\ d_2 \\ \vdots \\ d_{k-1} \end{bmatrix}$$

Similarly

$$Q^T A Q = \left. \begin{array}{c|c} 0 & 0 \\ \hline 0 & B \end{array} \right\} \quad k-1$$

$k-1$

Hence, (27) really takes the form

$$\left[\begin{array}{c|c} 0 & 0 \\ \hline 0 & B \end{array} \right] \begin{bmatrix} w \\ z_1 \\ z_2 \\ \vdots \\ z_{k-1} \end{bmatrix} = \begin{bmatrix} 0 \\ d_1 \\ d_2 \\ \vdots \\ d_{k-1} \end{bmatrix} \quad (28)$$

Clearly w is arbitrary in (28). The equation holds whenever

$$Bz = d \quad (29)$$

a $(k-1)$ -dimensional problem. All solutions of (26) are given by

$$x = Q \begin{bmatrix} w \\ z \end{bmatrix}$$

where z solves (29). The unique minimal length solution is given by

$$x^* = Q \begin{bmatrix} 0 \\ z \end{bmatrix}$$

By the assumptions on A , the matrix B is positive-definite and the conjugate gradient iteration can be used to solve (29). The condition number

$$\kappa(B) = \frac{\lambda_{k-1}(B)}{\lambda_1(B)} = \frac{\lambda_k(A)}{\lambda_2(A)}$$

(where we assume the eigenvalues are ordered algebraically) governs the convergence of the algorithm.

We now show that we can solve the consistent linear system (26) directly by the conjugate gradient algorithm, with convergence rate determined by $\kappa(B)$. The proof is elementary: we show that the algorithm, when started with an initial vector x_1 and a right hand side b which lie in the same invariant subspace of A , solves the problem while remaining entirely in that subspace.

Suppose we have the equation

$$Ax = b$$

with $b^T n = 0$ and $x_1^T n = 0$. Then the conjugate gradient algorithm, when applied to A , produces vectors x_i , r_i and p_i , exactly equal to

$$x_i = Q \begin{bmatrix} 0 \\ \vdots \\ z_i \end{bmatrix}$$

$$r_i = Q \begin{bmatrix} 0 \\ \vdots \\ \hat{r}_i \end{bmatrix}$$

and

$$p_i = Q \begin{bmatrix} 0 \\ \vdots \\ \hat{p}_i \end{bmatrix}$$

where z_i , \hat{r}_i and \hat{p}_i are the vectors produced by the conjugate gradient algorithm for

$$Bz = d$$

with initial vector $z_1 =$ the last $k-1$ entries of $Q^T x_1$. We note that

$$Q^T r_i = Q^T (b - Ax_i) = Q^T b - Q^T A (QQ^T) x_i = \begin{bmatrix} 0 \\ \vdots \\ d \end{bmatrix} - \begin{bmatrix} 0 \\ \vdots \\ Bz_i \end{bmatrix} = \begin{bmatrix} 0 \\ \vdots \\ \hat{r}_i \end{bmatrix},$$

where it is essential that $b^T n = 0$. It follows that

$$Q^T p_i = \begin{bmatrix} 0 \\ \vdots \\ \hat{p}_i \end{bmatrix}.$$

So assume the required properties hold for x_i , r_i and p_i . We shall show the iteration preserves these properties for x_{i+1} , r_{i+1} , and p_{i+1} . The first step in the iteration is to compute

$$\alpha_i = \frac{r_i^T r_i}{p_i^T A p_i},$$

which by induction is

$$\alpha_i = \frac{\begin{bmatrix} 0 \\ \vdots \\ \hat{r}_i \end{bmatrix}^T \begin{bmatrix} 0 \\ \vdots \\ \hat{r}_i \end{bmatrix}}{\begin{bmatrix} 0 \\ \vdots \\ \hat{p}_i \end{bmatrix}^T \begin{bmatrix} 0 \\ \vdots \\ \hat{p}_i \end{bmatrix}}$$

$$= \frac{\hat{r}_i^T \hat{r}_i}{\hat{p}_i^T B \hat{p}_i},$$

since $Q^T Q = I$ and

$$Q^T A Q = \begin{bmatrix} 0 & & 0 \\ & \text{---} & \\ 0 & & B \end{bmatrix}$$

Hence, $\hat{\alpha}_i = \alpha_i$, and it follows that

$$x_{i+1} = x_i + \alpha_i p_i = Q \begin{bmatrix} 0 \\ z_i \end{bmatrix} + \alpha_i Q \begin{bmatrix} 0 \\ \hat{p}_i \end{bmatrix} = Q \begin{bmatrix} 0 \\ z_i + \hat{\alpha}_i p_i \end{bmatrix} = Q \begin{bmatrix} 0 \\ z_{i+1} \end{bmatrix}$$

Then

$$\begin{aligned} r_{i+1} &= r_i - \alpha_i A p_i \\ &= Q \begin{bmatrix} 0 \\ \hat{r}_i \end{bmatrix} - \alpha_i A Q \begin{bmatrix} 0 \\ \hat{p}_i \end{bmatrix} \\ &= Q \begin{bmatrix} 0 \\ \hat{r}_i \end{bmatrix} - \alpha_i Q (Q^T A Q) \begin{bmatrix} 0 \\ \hat{p}_i \end{bmatrix} \\ &= Q \begin{bmatrix} 0 \\ \hat{r}_i - \hat{\alpha}_i B \hat{p}_i \end{bmatrix} = Q \begin{bmatrix} 0 \\ \hat{r}_{i+1} \end{bmatrix} \end{aligned}$$

That $\hat{\beta}_i = \beta_i$ follows from exactly the same reasoning that showed equality for the numerators of α_i and $\hat{\alpha}_i$. It is equally easy to show that

$$p_{i+1} = r_{i+1} + \beta_i p_i \iff Q^T p_{i+1} = \begin{bmatrix} 0 \\ \hat{r}_{i+1} + \hat{\beta}_i \hat{p}_i \end{bmatrix} = \begin{bmatrix} 0 \\ \hat{p}_{i+1} \end{bmatrix}.$$

Thus the conjugate gradient iteration will succeed in solving a consistent system.

We should note that the condition that $x_i^T n = 0$ is required only to assure that the solution also satisfies this property. If the minimal length solution is not desired, this condition can be ignored. The condition that $b^T n = 0$, that $Ax = b$ be consistent, is essential. If the system is not consistent, the initial residual r_i , and all successive residuals, will have the same component in the direction of the nullspace. Suppose that b (and r_i and p_i) has a component $\gamma \cdot n$ of the nullspace. The recursion for $\{r_i\}$

$$r_{i+1} = r_i - \alpha_i A p_i$$

shows that r_{i+1} has a component $\gamma \cdot n$ of the nullspace for all i . This, of course, also follows from the property that r_{i+1} is a residual vector for Ax_i and b . The directions $\{p_i\}$, and the approximate solutions $\{x_i\}$, have components of n which increase with i (in fact, rapidly).

It suffices to examine the recursion for p_i

$$p_{i+1} = r_{i+1} + \beta_i p_i$$

where $\beta_i > 0$. If δ_i is the component of n in the vector p_i ,

$$\delta_{i+1} = (\gamma + \beta_i \delta_i) = (\gamma + \beta_i (\gamma + \beta_{i-1} \delta_{i-1})) = \dots$$

Thus the direction's component of the nullspace increases. At the same time, p_{i+1} is shorter than r_{i+1} in the A norm, so the relative component in the direction of the nullspace grows. In the limit, the directions may become nearly parallel (they are still A -orthogonal, but A does not induce a metric). We can see this also by looking at the limiting case when an iteration is started with an actual solution vector. The residual vector (which is not a descent direction) and the initial direction, are in the nullspace—the initial step α_1 is infinite.

The last example given above implies that the convergence rate for obtaining the least squares solution to

an inconsistent system no longer depends on $\kappa(B)$. The convergence obtained in practice may be much slower than the bound obtainable for a consistent system. In practice, we are interested only in minimizing the residual, which we can ensure by working with a consistent system. In actual finite precision computing, the residuals and directions may wander slightly into the nullspace, in which case we are in the same situation as if we started with a slightly inconsistent system. The effects will be negligible unless the rate of convergence is very slow. In any case, the effects may be suppressed by explicitly reorthogonalizing the directions p_i to the nullspace.

In the preconditioned conjugate gradient algorithm we replace the system

$$Ax = b$$

by

$$Cz = d$$

where

$$C = VAV,$$

V some positive definite symmetric matrix. In the special case of a semi-definite A , a natural choice for a scaling matrix may also be semi-definite. This is the case for the two operators discussed in section 3. We have two possible cases to consider: V positive definite and V semi-definite.

The case of a positive definite V is little changed from the ordinary semi-definite conjugate gradient algorithm. We take $M = (V^2)^{-1}$, or more directly, $V = M^{-1/2}$. The rank of the matrix C is $k-1$ and its nullspace is

$$\text{span}\{M^{1/2}n_A\}$$

where n_A denotes a non-zero vector in the nullspace of A . The convergence condition for C required that \bar{r}_i and \bar{p}_i be orthogonal to $M^{1/2}n_A$.

The former condition will hold if $Ax = b$ is consistent since then

$$\bar{r}_i = M^{-1/2}(b - Ax_i)$$

and

$$(b - Ax_i)^T n_A = 0$$

The latter condition is equivalent to

$$p_i \perp Mn_A$$

since

$$p_i = M^{-1/2} \bar{p}_i$$

But clearly $p_i = M^{-1} r_i$ is orthogonal to Mn_A , since $r_i^T n_A = 0$. By induction, if $p_i^T Mn_A = 0$,

$$p_{i+1}^T Mn_A = (r_{i+1}^T M^{-1}) Mn_A + \beta_i (p_i^T Mn_A) = r_{i+1}^T n_A = 0.$$

Thus, convergence is governed by the pseudo-condition number $\frac{\lambda_{\kappa}(C)}{\lambda_{\lambda}(C)}$, the residual vectors r_i remain those of a consistent system, and the coefficients α_i and β_i are determined by an iteration remaining in a subspace

of the range of C . However, the scaled directions p_i are allowed to venture into the nullspace of A , and hence, the solution vector x is no longer necessarily the minimal length solution. This may be repaired easily at the end of the iteration by computing

$$x^* = x - \frac{x^T n_A}{n_A^T n_A} n_A.$$

The other natural case is one in which the scaling matrix M is also positive semi-definite, with the nullspace generated by a known vector n_M . The linear operator corresponding to M^{-1} in the definite case will be M^* , the pseudo-inverse of M . (For our purposes, we will need only a method for solving consistent systems $Mz = w$. Our knowledge of the nullspace of M then enables us to compute the minimal length least squares solution, M^*u , for any right hand side u). Formally, the matrix V appearing in the definition of C is $(M^*)^{1/2}$, so

$$C = (M^*)^{1/2} A (M^*)^{1/2}. \quad (6)$$

Several conditions which must be met by M are immediate. Since $\text{nullspace}(C) \supseteq \text{nullspace}(M)$, the rank of M must be at least the rank of A . For our specific problem $\text{rank}(M)$ must be at least $k-1$. Otherwise C has rank at most $k-2$ and cannot possibly span the entire $k-1$ dimensional space of possible solutions to $Ax = b$. To guarantee that C has rank $k-1$, we must have $\text{rank}(M) \geq k-1$ and also that the nullspace of M is not orthogonal to the nullspace of A . If the latter condition fails, we know that n_A lies in the range of M (since it is orthogonal to the orthogonal complement of the range of M). Hence, by symmetry of M , n_A lies in the range of $(M^*)^{1/2}$. There exists a vector z such that $(M^*)^{1/2} z = n_A$ and $z \perp n_M$. We would have:

$$C n_M = (M^*)^{1/2} A ((M^*)^{1/2} n_M) = (M^*)^{1/2} A 0 = 0$$

$$C z = (M^*)^{1/2} A ((M^*)^{1/2} z) = (M^*)^{1/2} A n_A = (M^*)^{1/2} 0 = 0$$

Thus, the dimension of the nullspace of C would be at least two.

Recall now the motivation for the preconditioned algorithm. The rate of convergence of the conjugate gradient algorithm is determined by the condition number, $\kappa(A)$, which is large for most discrete elliptic operators. To accelerate the coverage of the conjugate gradient algorithm, we replaced A by the preconditioned operator $C = M^{-1/2} A M^{-1/2}$, where we presume that $\kappa(C) \approx 1$. Our presumption can only be true if $\kappa(A) \approx \kappa(M)$, i.e., the preconditioning matrix must have roughly the same poor conditioning as has A . This requirement follows from the inequality

$$\kappa(C) \geq \max \left\{ \frac{\kappa(A)}{\kappa(M^{-1})}, \frac{\kappa(M^{-1})}{\kappa(A)} \right\}$$

where we note that $\kappa(M^{-1}) = \kappa(M)$. A brief proof of this inequality is given below:

$$\kappa(C) = \lambda_{\max}(C) / \lambda_{\min}(C).$$

However, the eigenvalues of C are the same as the eigenvalues of $M^{-1} A$. We can bound the eigenvalues of this latter product as:

$$\lambda_{\max}(M^{-1} A) \geq \lambda_{\max}(M^{-1}) \cdot \lambda_{\min}(A),$$

$$\lambda_{\min}(M^{-1} A) \leq \lambda_{\min}(M^{-1}) \cdot \lambda_{\max}(A).$$

It follows that

$$\begin{aligned}\kappa(C) &\geq \frac{\lambda_{\min}(A)}{\lambda_{\max}(A)} \cdot \frac{\lambda_{\max}(M)}{\lambda_{\min}(M)} \\ &= \kappa(M) / \kappa(A).\end{aligned}$$

The other inequality above follows by a similar argument.

We can obtain suitable preconditioning matrices for most positive definite discrete elliptic operators only by using poorly conditioned matrices. The required poor conditioning of M may be even worse in the case where both A and M are semi-definite. The preconditioning breaks down if

$$n_A^T n_M = 0.$$

We shall now show that the preconditioning nearly breaks down, in the sense that M must be very ill-conditioned, if the nullspaces are almost orthogonal. Assume that n_A and n_M have length one, and that

$$n_A^T n_M = \delta,$$

which is much less than one. We can write A and M^* in their respective eigendecompositions

$$A = U D_A U^T,$$

$$M^* = V D_{M^*} V^T,$$

where the zero eigenvalue of A and of M^* is given first. Then

$$C = M^{*1/2} A M^{*1/2} = S S^T,$$

where S is defined by

$$S = V D_{M^*}^{1/2} V^T U D_A^{1/2}.$$

For each of the matrices S , M^* , A , and X , define the condition number of the matrix as the condition number of the matrix restricted to the orthogonal complement of its nullspace. Then

$$\kappa(S) = \kappa[(D_{M^*}^{1/2} X) (D_A^{1/2})].$$

In the above equation, the matrix X is the $(k-1)$ by $(k-1)$ matrix obtained by removing the first row and column from $V^T U$, that is,

$$V^T U = \begin{bmatrix} \delta & a^T \\ b & X \end{bmatrix}$$

We now note that the inequality

$$\kappa(BC) \geq \max \left\{ \frac{\kappa(B)}{\kappa(C)}, \frac{\kappa(C)}{\kappa(B)} \right\}$$

holds for unsymmetric matrices B and C . (The proof follows from the triangle inequality for matrix norms,

$$\|B\| \leq \|BC\| \|C^{-1}\|.)$$

When we extend this inequality to products of three matrices we obtain the relatively weak result that

$$\kappa(BCD) \geq \max \left\{ \frac{\kappa(B)}{\kappa(C)\kappa(D)}, \frac{\kappa(C)}{\kappa(B)\kappa(D)}, \frac{\kappa(D)}{\kappa(B)\kappa(C)} \right\}$$

In the context of our preconditioned operator, only $\kappa(X)$ is unknown. However, Alan Cline [11] has shown that when $\delta \ll 1$,

$$\kappa(X) \cong 1 / \delta.$$

Thus, the condition number of the preconditioned operator C is bounded below by:

$$\kappa(C) \geq (1 / \delta)^2 / (\kappa(A) \kappa(M)).$$

Since $\kappa(A)$ is fixed, this implies that $\kappa(M)$ must be large whenever δ is small.

V. Numerical Results

A code, written in FORTRAN, was developed to implement the preconditioned conjugate gradients scheme discussed in previous sections. Care was taken in the preparation of the code to make it portable and to introduce many comments for clarity. The code has been run successfully under a variety of conditions and has been compared with analytical results to determine its accuracy. In this section a description is given of some of the computations used to determine the performance of this code.

The model problem, the pressure equation, was discussed in section III. Special cases of this general non-separable elliptic equation were used to test the code for accuracy. All production runs of this code were performed when the code was imbedded in a larger linear or nonlinear fluid dynamics computation; timing studies were performed in such an environment.

Within the code, two tests are used to terminate the conjugate gradient iteration; these depend upon two specified parameters, the maximum number of iterations (less than or equal to 50) and a maximum relative residual, ϵ . (The relative residual is defined as the norm of the residual divided by the norm of the right hand side in the scaled problem discussed in sections 2 and 4.) In all successful computations to date, the iteration is terminated after a relatively small number of iterations by the relative residual norm satisfying the criterion that it be below ϵ .

To test the code under the simplest conditions, a Poisson equation on the unit square was discretized and solved with homogeneous Neumann conditions applied at the boundary. In continuous form this problem is

$$\frac{\partial^2 p}{\partial x^2} + \frac{\partial^2 p}{\partial y^2} = \tilde{f}(x, y) \equiv \cos(x\pi) \cos(l\pi y)$$

on $0 \leq x \leq 1$ and $0 \leq y \leq 1$. The boundary conditions are

$$\frac{\partial p}{\partial x} = 0 \text{ at } x = 0 \text{ and } x = 1,$$

$$\frac{\partial p}{\partial y} = 0 \text{ at } y = 0 \text{ and } y = 1.$$

The solution to this problem is

$$p(x, y) = - \frac{1}{(l\pi)^2 + (x\pi)^2} \cos(x\pi) \cos(l\pi y)$$

When this problem is discretized to second order accuracy on the "staggered grid" (a grid displaced one half incremental unit δx in the x -direction and one half incremental unit δy in the y -direction) as discussed in section III, it can be written

$$\frac{1}{\delta x^2} (\bar{p}_{i+1,j} - 2\bar{p}_{i,j} + \bar{p}_{i-1,j}) + \frac{1}{\delta y^2} (\bar{p}_{i,j+1} - 2\bar{p}_{i,j} + \bar{p}_{i,j-1}) = \bar{f}_{ij}$$

Where

$$\bar{f}_{ij} = \cos\left[\frac{x\pi}{m}(i-1/2)\right] \cos\left[\frac{y\pi}{n}(j-1/2)\right]$$

for

$$1 \leq i \leq m \text{ and } 1 \leq j \leq n.$$

Here δx is the mesh spacing in the x -direction, $\delta x = 1/m$ (m is the number of mesh cells in the x -direction) and δy is the mesh spacing in the y -direction, $\delta y = 1/n$ (n is the number of cells in the y -direction). The boundary conditions are

$$\bar{p}_{0,j} = \bar{p}_{1,j} \text{ and } \bar{p}_{m+1,j} = \bar{p}_{m,j} \text{ for } 1 \leq j \leq n$$

$$\bar{p}_{i,0} = \bar{p}_{i,1} \text{ and } \bar{p}_{i,n+1} = \bar{p}_{i,n} \text{ for } 1 \leq i \leq m$$

The solution to this linear algebraic system is

$$\bar{p}_{ij} = \frac{1}{\left(2m \sin \frac{x\pi}{2m}\right)^2 + \left(2n \sin \frac{y\pi}{2n}\right)^2} \cos\left[\frac{x\pi}{m}(i-1/2)\right] \cos\left[\frac{y\pi}{n}(j-1/2)\right]$$

for $0 \leq i \leq m+1, 0 \leq j \leq n+1$

The solution to this simple discretized Poisson equation was computed using the code, and the solution compared with the analytical solution given above.

A small test problem, $m = n = 7$, was used to determine the accuracy obtainable when $\epsilon = 10^{-6}$. Comparison with the exact solution demonstrated that the components of the solution vector obtained from the computation agreed to at least six significant figures with the exact solution. Generally the agreement was much better, being seven or eight significant figures for most values of i and j . (Note that the UNIVAC 1108, on which all computations were run, carries about eight significant figures.)

As a larger test problem, the equations with $m = n = 31$ were solved with $\epsilon = 10^{-6}$. For this computation agreement was obtained to a few parts in the sixth significant figure.

A second test problem, a discretized approximation to a separable elliptic equation, was also solved analytically and using the code. Comparison of these results indicated that the accuracy was similar to that obtained in the first test problem.

The code was then imbedded in a linear finite difference computation obtained from a second-order discretization of a set of equations arising in fluid dynamics. These equations describe two-dimensional internal gravity waves in a stratified ambient fluid within a rectangular enclosure. The interest in this problem is discussed, the continuous and discrete equations are presented and exact analytical solution to the continuous and discrete problem are given by Baum and Rehm [8]. Additional linear computations using this code on a somewhat more general fluid-flow problem are described in Rehm and Baum [7]. In this latter paper the more general linear finite difference equations are given, and the computational procedure for solving them is presented. The manner in which the preconditioned conjugate gradients code is used in solving the elliptic pressure equation is discussed. In all of these linear computations, the equation for the pressure is separable; it is only in the general, nonlinear computations that the equation for the pressure is nonseparable.

For the linear computations reported in these papers, the number of iterations required to obtain a relative residual error less than $\epsilon = 10^{-5}$ or 10^{-6} generally varied between 2 and 4. A large number of such computations have been run.

A representative one is a calculation for which $m = 15$, $n = 16$ and $\epsilon = 10^{-6}$; in this computation the pressure-solver was called 200 times. Two iterations to convergence were taken ten of the first eleven times it

was called, each call using about 0.5 s of CPU time on the NBS Univac 1108. Thereafter, each subsequent call required only one iteration to convergence taking about 0.35 s of CPU. An indication of the rate of convergence of the algorithm is obtained by taking

$$\delta \equiv \frac{\log_{10}(\text{Initial relative residual}) - \log_{10}(\text{final relative residual})}{\text{number of iterations}}$$

For the computation reported above, this rate of convergence was about 4.5 when two iterations were taken and about 5.5 when one iteration was taken.

It should be noted that these computations determine a flow field which evolves with time. Except for the first time step, the pressure vector at the previous time step is used as the initial guess for the pressure vector each time the elliptic-solver is called. Hence the guess at each calling is quite good and convergence is rapid.

Computations of a set of nonlinear finite difference equations generalizing the linear ones described above have also been run. As noted previously, in this case the elliptic equation for the pressure is nonseparable. In a limited number of computations, the PCG code has been found to perform well in this case also. A representative computation, one for which $I = 31, J = 31$ and $\epsilon = 10^{-5}$, was found to take between 2 and 5 iterations for convergence at each call. The CPU time taken was slightly over 2 s for 2 iterations and slightly under 5 s for 5 iterations (very roughly a second per iteration generally). In this calculation δ defined above was found to vary between about one and two.

In the Appendix to [15] a listing of the elliptic-solver, called FASTSL, is given. To use this code, subroutines from EISPACK, the Argonne Code Center package, and BLKTRI, a subroutine in the NCAR package developed by Schwarztrauber and Sweet [10] are required.

The authors wish to acknowledge the help of our colleagues, Mr. Martin Cordes of the Center for Applied Mathematics (C.A.M.) and Dr. Howard Baum of the Center for Fire Research (C.F.R.). This study would not have been completed without their help. We also wish to thank Drs. Dianne O'Leary and Francis Sullivan for reading and correcting the manuscript. One of us (J. L.) would like to express his appreciation to Drs. B. H. Colvin and F. C. Johnson for providing the opportunity to undertake this study in the C.A.M. during the summer of 1977.

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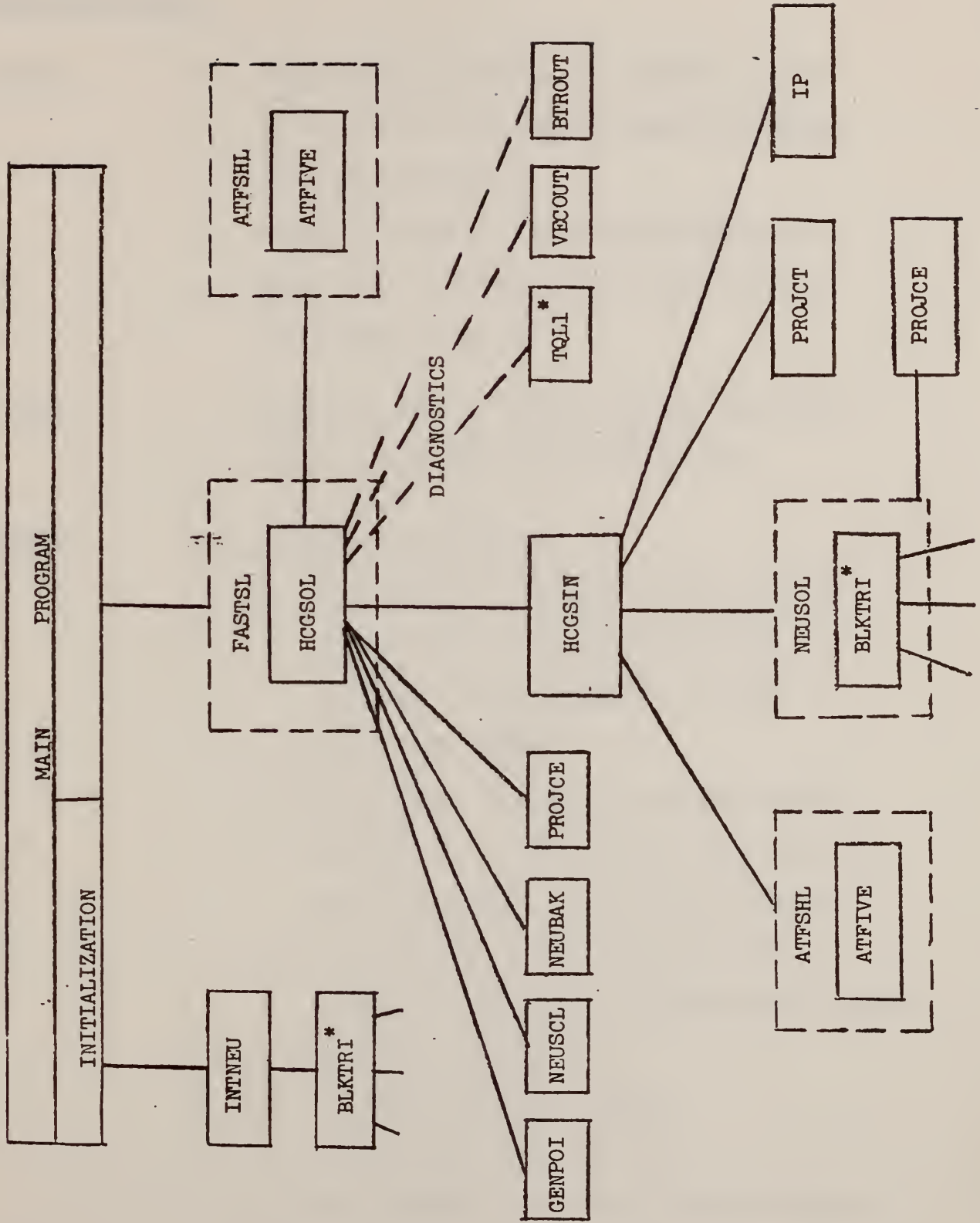
APPENDIX

DIAGRAM OF SUBROUTINE LINKAGES

BRIEF DESCRIPTION OF SUBROUTINES

LISTING OF ALL SUBROUTINES

DIAGRAM OF SUBROUTINE LINKAGES



* Subroutines from external libraries

Brief Description

- INTNEU - Initializes the NCAR Routine BLKTTRI to solve $\nabla^2 u = f$ on the rectangular grid (with Neumann Boundary Conditions)
Generates parameter vector PPARAM which must be passed back to BLKTTRI when we actually solve the pressure equation
- FASTSL - Separates required working storage into pieces for the coordinating subroutine HCGSOL
- HCGSOL - Coordinates the pieces
- a) Calls GENPOI to generate the discrete five point operator for the pressure equation.
The matrix's representation is stored in the vector 'APARM'
 - b) Calls NEUSCL to compute a diagonal scaling to make the discrete pressure equation have the same diagonal as the Poisson equation.
 - c) Calls HCGSIN with the diagonally scaled operator. HCGSIN uses preconditioned conjugate gradient algorithm with
A \equiv scaled pressure equation
P \equiv discrete Poisson equation
 - d) Calls NEUBAK to undo the diagonal scaling to get a solution of the original pressure equation

The first part of the report deals with the general situation of the country in 1900. It is a year of transition, a year of change. The old is passing away, and the new is coming. The people are restless, they are looking for something better, something more. They are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The second part of the report deals with the economic situation. It is a year of economic change, a year of economic transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The third part of the report deals with the social situation. It is a year of social change, a year of social transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The fourth part of the report deals with the political situation. It is a year of political change, a year of political transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The fifth part of the report deals with the cultural situation. It is a year of cultural change, a year of cultural transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The sixth part of the report deals with the educational situation. It is a year of educational change, a year of educational transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The seventh part of the report deals with the religious situation. It is a year of religious change, a year of religious transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The eighth part of the report deals with the military situation. It is a year of military change, a year of military transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The ninth part of the report deals with the foreign situation. It is a year of foreign change, a year of foreign transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

The tenth part of the report deals with the future. It is a year of future change, a year of future transition. The old is passing away, and the new is coming. The people are tired of the old ways, they are tired of the old life. They want to know what the future holds for them, and they want to know what they can do to make it better.

HCGSOL also coordinates diagnostic information

HCGSIN - Preconditioned conjugate gradients repeated by
 calls to subroutine to compute
 Ax for some x
 and a subroutine to solve $Py = x$
 for some x

The names of the subroutines are parameters from
HCGSOL. Here they are

ATF5HL and NEUSOL

ATF5HL - Unpacks representation of scaled pressure equation
 in APARM to pass it to ATFIVE, which actually per-
 forms the matrix-vector product.

NEUSOL - Unpacks representation of POISSON equation in PPARM,
 to pass it to BLKTRI to actually use cyclic re-
 duction to solve the discrete Poisson equation.

Other Subroutines

PROJCT, PROJCE - To compute the projection of an arbitrary vector
 on the orthogonal complement of some vector
 (PROJCT) or the vector of all one's (PROJCE)

IP - To compute the inner product of two vectors

MEMORANDUM FOR THE RECORD

1. On 10/10/50, the following information was received from the

Department of the Interior:

2. The following information was received from the

Department of the Interior:

3. The following information was received from the

Department of the Interior:

4. The following information was received from the

Department of the Interior:

5. The following information was received from the

Department of the Interior:

6. The following information was received from the

Department of the Interior:

7. The following information was received from the

Department of the Interior:

8. The following information was received from the

Department of the Interior:

9. The following information was received from the

Department of the Interior:

BTROUT - Prints out a symmetric block matrix with five non-zero diagonals used only for diagnostic output.

VECOUT - Prints a vector in a grid-like form. Used only for diagnostic output.

Subroutines from external libraries

TQL1 - From EISPACK. Computes all eigenvalues of tri-diagonal matrix, (produced implicitly by conjugate gradients); used to estimate convergence rate of conjugate gradients

BLKTRI - NCAR routine. Used to solve $\nabla^2 P = f$ with Neumann boundary conditions.

***** ATFIVE *****

@:ELT.L -PF1-.ATFIVE
ELT BR1 S74Q1C 07/11/79 08:48:30 (1)

SUBROUTINE ATFIVE (M, N, D1, D2, DM, X, AX)

=====
MULTIPLY X BY A WHERE A IS THE REPRESENTATION
OF A SPARSE FIVE POINT OPERATOR
=====

INTEGER M, N
REAL X(1), AX(1), D1(1), D2(1), DM(1)

INTEGER MP1, NMM1, NMEM, NM, I, NMMMP1
NM = N * M

... 1ST BLOCK ...
AX(1) = D1(1) * X(1) + D2(1) * X(2) + DM(1) * X(M+1)
DO 100 I = 2, M
AX(I) = D1(I) * X(I) + D2(I) * X(I+1) + D2(I-1) * X(I-1)
+ DM(I) * X(I+M)
100 CONTINUE

... INTERIOR BLOCKS ...

MP1 = M + 1
NMEM = NM - M
DO 200 I = MP1, NMM1
AX(I) = D1(I) * X(I) + DM(I-M) * X(I-M) + D2(I-1) * X(I-1)
+ DM(I) * X(I+M) + D2(I) * X(I+1)
200 CONTINUE

NMMMP1 = NMM1 + 1
NMM1 = NM - 1

... FINAL BLOCK ...

DO 300 I = NMMMP1, NMM1
AX(I) = D1(I) * X(I) + DM(I-M) * X(I-M) + D2(I-1) * X(I-1)
+ D2(I) * X(I+1)
300 CONTINUE
AX(NM) = D1(NM) * X(NM) + D2(NMM1) * X(NMM1)
+ DM(NMEM) * X(NMEM)

RETURN
END

END ELT. ERRORS: NONE. TIME: 0.209 SEC. IMAGE COUNT: 49

@HDC.P ***** ATFSHL *****

@:ELT.L -PF1-ATFSHL

ELT BR1 S74QIC 07/11/79 08:48:32 (3)

1. 03 SUBROUTINE ATFSHL (APARM, X, AX)

2. 00

3. 00

4. 00

5. 02

6. 00

7. 00

8. 03

9. 00

10. 03

11. 00

12. 00

13. 00

14. 02

15. 00

16. 03

17. 03

18. 00

19. 00

20. 00

21. 00

22. 00

23. 00

24. 03

25. 00

26. 00

27. 00

28. 00

=====
SHELL ROUTINE FOR MATRIX-VECTOR MULTIPLY FOR DISCRETE
FIVE POINT OPERATOR ON A RECTANGULAR GRID.
=====

REAL APARM(1), X(1), AX(1)

VECTOR APARM CONTAINS PARAMETERS FOR MULTIPLY

INTEGER M, N, NM, SUBD1, SUBD2, SUBD3

DECODE PARAMETERS AND PASS ON TO 'ATFIVE'

M = FIX(APARM(1))

N = FIX(APARM(2))

NM = N * M

SUBD1 = 3

SUBD2 = 3 + NM

SUBD3 = SUBD2 + NM

CALL ATFIVE (M, N, APARM(SUBD1), APARM(SUBD2), APARM(SUBD3),
1 X, AX)

RETURN

END

END ELT. ERRORS: NONE. TIME: 0.159 SEC. IMAGE COUNT: 28

@HDC,P ***** BTROUT *****

@:ELT,L -PF1-.BTROUT

ELT 8R1 S74Q1C 07/11/79 08:48:33 (4)

SUBROUTINE BTROUT (M, N, D1, D2, D3)

```

1. 00 C
2. 00 C
3. 00 C
4. 00 C
5. 00 C
6. 00 C
7. 00 C
8. 02 C
9. 00 C
10. 00 C
11. 00 C
12. 00 C
13. 00 C
14. 00 C
15. 00 C
16. 01 C
17. 00 C
18. 00 C
19. 02 C
20. 04 C
21. 04 C
22. 04 C
23. 01 C
24. 04 C
25. 02 C
26. 04 C
27. 04 C
28. 04 C
29. 04 C
30. 00 C
31. 00 C
32. 00 C
33. 00 C
34. 00 C
35. 00 C
36. 00 C
37. 00 C
38. 00 C
39. 00 C

```

 PRINT THE UPPER TRIANGLE OF A BLOCK TRIDIAGONAL MATRIX
 RESULTING FROM A FINITE DIFFERENCE APPROXIMATION ON A
 RECTANGULAR GRID USING A FIVE POINT OPERATOR.

M : BLOCK SIZE
 N : NUMBER OF BLOCKS
 D1 : MAIN DIAGONAL
 D2 : 1ST SUPERDIAGONAL
 D3 : MAIN DIAGONAL OF OFF-DIAGONAL BLOCKS

INTEGER M, N
 D1(1), D2(1), D3(1)
 ACTUAL DIMENSION IS M * N

INTEGER I, J, I1, I2

INTEGER PRINTR
 DATA PRINTR / 6 /

WRITE (PRINTR, 60000)

DO 10 J = 1, N

I1 = (J-1) * M + 1

I2 = J * M

WRITE (PRINTR, 60100) (D1(I), D2(I), D3(I), I = I1, I2)

WRITE (PRINTR, 60200)

10 CONTINUE

RETURN

 60000 FORMAT ('OUTPUT BY R O W S OF MATRIX'/)
 60100 FORMAT (1P2E15.7, 10X, E15.7)
 60200 FORMAT (' ')

END ELT. ERRORS: NONE. TIME: 0.201 SEC. IMAGE COUNT: 39

@HDC,P ***** DIRBAK *****

***** FASTSL *****

0:ELT,L -PF1-,FASTSL
 ELT BR1 5740IC 07/11/79 08:48:56 (15)

1 11 SUBROUTINE FASTSL (M, N, MMAX, RHO, PRESSR, RES, RESID,
 2 12 LEFTB, RIGHTB, LOWERB, UPPERB, PPARM,
 3 12 DELTAX, DELTAY, ITERS, RTOL,
 4 12 WORK, CNTL)
 5 11
 6 11 M, N, MMAX
 7 12 RHO(MMAX, N), PRESSR(MMAX, N), RES(MMAX, N),
 8 12 RESID(MMAX, N)
 9 12 LEFTB(N), RIGHTB(N), LOWERB(M), UPPERB(M), PPARM(1)
 10 12 DELTAX, DELTAY, RTOL
 11 12 ITERS
 12 11 WORK(1)
 13 11 CNTL
 14 11
 15 11
 16 11
 17 11
 18 11
 19 11
 20 11
 21 11
 22 12
 23 12
 24 11
 25 11
 26 11
 27 11
 28 11
 29 11
 30 12
 31 11
 32 12
 33 11
 34 12
 35 12
 36 12
 37 12
 38 11
 39 11
 40 11
 41 11
 42 11
 43 11
 44 11
 45 12
 46 12
 47 12
 48 12
 49 12
 50 12
 51 12
 52 12
 53 12
 54 12
 55 12
 56 12

 SOLVE THE PRESSURE EQUATION WITH NEUMANN BOUNDARY CONDITIONS ON
 A RECTANGLE, USING HYBRID CONJUGATE GRADIENTS WITH CYCLIC
 REDUCTION AS THE FAST METHOD INSIDE THE ITERATION.

PARAMETERS:
 (STARRED PARAMETERS ARE ALTERED BY THE SUBROUTINE)

- M - THE NUMBER OF GRID POINTS IN THE HORIZONTAL DIRECTION
- N - THE NUMBER OF GRID POINTS IN THE VERTICAL DIRECTION
- MMAX - THE DECLARED ROW DIMENSION OF THE GRID ARRAYS
 (MUST BE AT LEAST M, MAY BE THE SAME AS M).
- RHO - THE DENSITIES AT THE GRID POINTS
- *PRESSR - AN INITIAL GUESS AT THE SOLUTION PRESSURE. THE FINAL
 SOLUTION IS RETURNED IN THIS ARRAY.
- *RES - THE RIGHT HAND SIDE OF THE ELLIPTIC EQUATION
 (SHOULD NOT INCORPORATE THE BOUNDARY CONDITIONS.)
 THE VALUES IN THE ARRAY WILL BE DESTROYED BY THIS
 SUBROUTINE.
- *RESID - THE RESIDUAL 'VECTOR' FOR THE LINEAR DIFFERENCE
 EQUATIONS EVALUATED AT THE FINAL SOLUTION 'PRESSR'.
 (THIS ARRAY DOES NOT NEED TO BE INITIALIZED.)
- LEFTB - THE NORMAL DERIVATIVE OF PRESSR, DIVIDED BY RHO,
 ON THE LEFT BOUNDARY.
- RIGHTB - THE BOUNDARY DERIVATIVES ON THE RIGHT BOUNDARY
- LOWERB - THE BOUNDARY DERIVATIVES ON THE LOWER BOUNDARY
- UPPERB - THE BOUNDARY DERIVATIVES ON THE UPPER BOUNDARY
- PPARM - A VECTOR OF PARAMETERS TO THE CYCLIC REDUCTION
 FAST POISSON SOLVER. THIS VECTOR IS INITIALIZED
 BY THE SUBROUTINE 'INTNEU', WHICH MUST BE CALLED
 (ONCE) BEFORE THIS FAST NON-SEPARABLE SOLVER IS
 CALLED. AS LONG AS THE CALLING PROGRAM DOES NOT
 ALTER THE CONTENTS OF PPARM, IT WILL NOT NEED
 TO BE REINITIALIZED.
 PPARM MUST BE OF LENGTH AT LEAST:
 $3N + 9M + 2(N+E)(\text{LOG}(N)-1) + 12$
 UNLESS
 N IS INAPPROPRIATE FOR THE FAST SOLVER

***** FASTSL *****

```

57. C
58. C
59. C
60. C
61. C
62. C
63. C
64. C
65. C
66. C
67. C
68. C
69. C
70. C
71. C
72. C
73. C
74. C
75. C
76. C
77. C
78. C
79. C
80. C
81. C
82. C
83. C
84. C
85. C
86. C
87. C
88. C
89. C
90. C
91. C
92. C
93. C
94. C
95. C
96. C
97. C
98. C
99. C
100. C
101. C
102. C
103. C
104. C
105. C
106. C
107. C
108. C
109. C
110. C
111. C
112. C
113. C
114. C

AND M IS APPROPRIATE
OR
BOTH ARE APPROPRIATE AND N IS THE LARGER,
IN WHICH CASE PPARM MUST BE AT LEAST AS LONG AS:
3M + 9N + 2(M+2)(LOG(M)-1) + 12.

DELTA - THE GRID SPACING IN THE HORIZONTAL DIRECTION
DELTA - THE GRID SPACING IN THE VERTICAL DIRECTION
*ITERS - ON INPUT, THE MAXIMUM NUMBER OF CG ITERATIONS ALLOWED
ON OUTPUT, THE NUMBER OF ITERATIONS ACTUALLY USED.
*RTOL - ON INPUT, THE RELATIVE RESIDUAL DEMANDED OF CONJUGATE
GRADIENTS (USING THE SCALED MATRIX). ON OUTPUT, THE
TWO NORM OF RESID. RELATIVE TO THE NORM OF THE
SOLUTION VECTOR.

*WORK - A VECTOR OF WORKING STORAGE, OF LENGTH >= 6NM + 2

*CNTRL - CONTROL FLAG. ON INPUT, CONTROLS THE AMOUNT OF PRINTING
PRODUCED BY THE FAST SOLVER.
VALID INPUT VALUES:
0 - SUPPRESS ALL OUTPUT
1 - SUMMARY OF CONJUGATE GRADIENT ITERATIONS
2 - ABOVE, PLUS SUMMARY STATISTICS ON INPUT AND
OUTPUT VECTORS.
3 - ABOVE, PLUS LISTING OF RMS, PRESSR AND RESID
4 - ABOVE, PLUS LISTING OF UNSCALED AND SCALED
MATRICES.

VALID OUTPUT VALUES:
0 - SUCCESSFUL SOLUTION
1 - FAILURE TO MEET RESIDUAL TOLERANCE
(SOME FATAL ERRORS CAUSE IMMEDIATE TERMINATION)

=====
LOCAL VARIABLES ...
INTEGER NM, SAFARM, SALPHA, SBETA, SCANMA, SNULL, SF, SX
EXTERNAL GENPOL, NEUBAK, NEUSCL, NEUSOL, AITFEHL
LOGICAL FLIP
INTEGER PRINTR
DATA PRINTR / 6 /

SPLIT THE WORK VECTOR INTO ARRAYS FOR CONJUGATE GRADIENTS
AND CALL THE COORDINATING SUBROUTINE HCGSOL.

NM = N * M
SAFARM = 1
SALPHA = SAFARM + 2
SBETA = SALPHA + NM
SCANMA = SBETA + NM
SNULL = SCANMA + NM
SF = SNULL + NM
SX = SF + NM

DECIDE WHICH ONE-DIMENSIONAL ORDERING WILL BE MOST

```


115. C
 116. C
 117. C
 118. C
 119. C
 120. C
 121. C
 122. C
 123. C
 124. C
 125. C
 126. C
 127. C
 128. C
 129. C
 130. C
 131. C
 132. C
 133. C
 134. C
 135. C
 136. C
 137. C
 138. C
 139. C
 140. C
 141. C
 142. C
 143. C
 144. C
 145. C
 146. C
 147. C
 148. C
 149. C
 150. C
 151. C
 152. C
 153. C
 154. C
 155. C
 156. C
 157. C
 158. C
 159. C
 160. C
 161. C
 162. C
 163. C
 164. C
 165. C
 166. C
 167. C
 168. C
 169. C
 170. C
 171. C
 172. C

APPROPRIATE FOR THE FAST SOLVER.

THE SECOND SUBSCRIPT MUST HAVE A LIMIT WHICH IS ONE LESS THAN A POWER OF TWO (FOR SUBROUTINE BLKTRI). SWITCH SUBSCRIPTS TO MEET THIS REQUIREMENT OR TO PUT THE SMALLER LIMIT SECOND IF BOTH ARE APPROPRIATE.

THE INITIALIZING SUBROUTINE FOR THE SCALING MATRIX MUST MAKE THE SAME DECISION ON THE ORDERING.

```

FLIP = ( ( N .NE. 2**(LOG2(N+1)) - 1) .OR.
1      ( ( N .EQ. 2**(LOG2(N+1)) - 1) .AND.
2      ( M .EQ. 2**(LOG2(M+1)) - 1) .AND.
3      ( M .LT. N ) )

```

CHECK DIMENSIONS OF SYSTEM -- MAKE SURE THAT THE FAST SOLVER HAS BEEN INITIALIZED WITH THE SAME DIMENSIONS AS ARE BEING USED HERE.

NB: IF THE STRUCTURE OF THE PARAMETER VECTOR 'PPARM' IS CHANGED (AS FOR A DIFFERENT FAST SOLVER), THIS CODE MAY BE AFFECTED.

```

IF ( (FLIP .AND.
1     (IFIX(PPARM(2)) .NE. N) .AND. (IFIX(PPARM(3)) .NE. M) )
2   .OR.
3   (.NOT. FLIP .AND.
4     (IFIX(PPARM(2)) .NE. M) .AND. (IFIX(PPARM(3)) .NE. N) )
5   .OR.
6   (PPARM(1) .EQ. 0.0) ) GO TO 1000

```

CALL HCGSOL (M, N, NM, NMAX, FLIP,

```

1 GENPOI, NEUSCL, NEUBAK, ATFSHL, NEUSOL,
2 RHO, RES, PRESSR, WORK(SF), WORK(SX), RESID,
3 RHS, PRESSR, RESID, WORK(SNELL), WORK(SAPARM),
4 WORK(SALPHA), WORK(SBETA), WORK(SCAMMA),
5 LOWERB, UPPERB, LEFTB, RIGHTB, PPARML,
6 DELTAY, DELTAY, ITERS, RTOL, CNTL)

```

(THE SECOND OCCURENCE OF PARAMETERS RES, PRESSR, AND RESID IN THE CALL ABOVE IS AS A WORKING ARRAY. ALTERATIONS TO THE LOGIC OF HCGSOL COULD REQUIRE ADDITIONAL TEMPORARY SPACE. NOTE THAT THESE ARRAYS ARE DESTROYED AS SOON AS THE ARRAYS WRK1, WRK2, WRK3 OF HCGSOL ARE USED.)

RETURN

ERROR HANDLING ...

```

1000 WRITE (PRINTR, 61000) M, N
STOP

```

```

61000 FORMAT ('*****
1 'THE FAST POISSON SOLVER HAS NOT BEEN INITIALIZED.'/
2 ' OR HAS BEEN INITIALIZED WITH DIFFERENT DIMENSIONS.'/
3 ' DIMENSIONS GIVEN TO FASTSL:.'/

```


***** FASTSL *****

DATE 071179

PAGE

173.	12	4	M: , 16 /
174.	12	5	N: , 16 /
175.	12	6	*****
176.	12		*****

END

END ELT. ERRORS: NONE. TIME: 0.651 SEC. IMAGE COUNT: 176

@HDG.P ***** GENCS3D *****

M * N. THE GRID POINTS ARE STORED IN EQUIVALENT ONE-DIMENSIONAL FASHION IN ONE OF TWO ORDERINGS. THE NATURAL (FORTRAN) ORDER IS SCANNING THE GRID ACROSS ROWS FROM LEFT TO RIGHT, MOVING FROM BOTTOM ROW TO TOP ROW. THIS ORDER WILL BE USED UNLESS IT IS MORE EFFICIENT (OR NECESSARY) TO USE THE TRANSPOSE OF THIS ORDER (UP COLUMNS, FROM LEFT TO RIGHT). FLIP IS FALSE UNLESS THE TRANSPOSED ORDER SHOULD BE USED.

WORKING STORAGE (TEMPORARY STORAGE)

DENS A VECTOR OF LENGTH AT LEAST M

P A R A M E T E R D E C L A R A T I O N S

INTEGER M, N, MMAX
 LOGICAL FLIP
 REAL RHO(MMAX,N), F(MMAX,N), PRESSR(MMAX,N)
 REAL DELTAX, DELTAY
 REAL LOWERB(YD), UPPERB(YD), LEFTS(N), RIGHTB(N)
 REAL ALPHA(1), BETA(1), GAMMA(1), FP(1), XP(1)
 REAL DENS(MMAX)

L O C A L V A R I A B L E S

INTEGER I, J, IJ
 REAL ODX2, ODY2, RDENS, LDENS, DDENS

I N I T I A L I Z E G R I D P A R A M E T E R S

ODX2 = 1.0 / DELTAX**2
 ODY2 = 1.0 / DELTAY**2

THE DENSITY FUNCTION IS APPROXIMATED AT THE MIDPOINTS OF GRID LINES, SINCE WE NEED APPROXIMATIONS TO THE DERIVATIVE OF THE DENSITY FUNCTION. THE VECTOR 'DENS' AND THE VARIABLE 'RDENS' ARE USED TO PASS THE MIDPOINT CALCULATIONS TO THE NEXT ROW AND THE NEXT COLUMN (TO AVOID RECALCULATION).

S T R U C T U R E

- 1) GENERATE VALUES FOR BOTTOM ROW
 - 2) GENERATE VALUES FOR INTERIOR ROWS
 - 3) GENERATE VALUES FOR TOP ROW
- STRUCTURE FOR EACH ROW:
- A) GENERATE VALUE FOR LEFT COLUMN
 - B) GENERATE VALUES FOR INTERIOR COLUMNS
 - C) GENERATE VALUE FOR RIGHT COLUMN

57. C
 58. C
 59. C
 60. C
 61. C
 62. C
 63. C
 64. C
 65. C
 66. C
 67. C
 68. C
 69. C
 70. C
 71. C
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 103. C
 104. C
 105. C
 106. C
 107. C
 108. C
 109. C
 110. C
 111. C
 112. C
 113. C
 114. C




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115. C 38
116. C 38
117. C 38
118. C 36
119. C 38
120. C 38
121. C 36
122. C 38
123. C 38
124. C 38
125. C 36
126. C 41
127. C 38
128. C 36
129. C 33
130. C 38
131. C 38
132. C 38
133. C 38
134. C 41
135. C 38
136. C 38
137. C 38
138. C 38
139. C 38
140. C 38
141. C 36
142. C 38
143. C 38
144. C 38
145. C 38
146. C 38
147. C 41
148. C 38
149. C 38
150. C 38
151. C 38
152. C 38
153. C 36
154. C 41
155. C 38
156. C 38
157. C 38
158. C 36
159. C 38
160. C 38
161. C 38
162. C 38
163. C 38
164. C 38
165. C 38
166. C 40
167. C 40
168. C 40
169. C 40
170. C 40
171. C 40
172. C 40

DO 200 J = 1, N
DO 100 I = 1, M

IF (.NOT. FLIP) IJ = (J-1)*M + I
IF (FLIP) IJ = (1-1)*N + J

IF ( J .CT. 1 ) GO TO 10
... FIRST ROW ...
DDENS = 0.0
UDENS(1) = 2.0 / ( RHO(1,1) + RHO(1,2) )
FP(IJ) = F(1,J) - LOWERD(1) / DELTAY
GO TO 30

IF ( J .EQ. N ) GO TO 20
... INTERIOR ROWS ...
DDENS = DENS(I)
UDENS = 2.0 / ( RHO(1,J) + RHO(1,J+1) )
FP(IJ) = F(1,J)
GO TO 30

... LAST ROW ...
DDENS = DENS(I)
UDENS = 0.0
FP(IJ) = F(1,J) - UPPERB(I) / DELTAY

IF ( I .CT. 1 ) GO TO 40
... FIRST COLUMN ...
LDENS = 0.0
RDENS = 2.0 / ( RHO(1,J) + RHO(2,J) )
FP(IJ) = FP(IJ) - LEFTB(J) / DELTAX
GO TO 60

IF ( I .EQ. M ) GO TO 50
... INTERIOR COLUMNS ...
LDENS = RDENS
RDENS = 2.0 / ( RHO(1,J) + RHO(1+1,J) )
GO TO 60

... LAST COLUMN ...
LDENS = RDENS
RDENS = 0.0
FP(IJ) = FP(IJ) - RIGHTB(J) / DELTAX

CALCULATE FIVE POINT OPERATOR AT POINT ( I, J )

ALPHA(IJ) = -ODX2 * (RDENS + LDENS)
            - ODY2 * (DDENS + UDENS)
IF (FLIP) GO TO 80
BETA(IJ) = ODX2 * RDENS
GAMMA(IJ) = ODY2 * UDENS
GO TO 90

BETA(IJ) = ODY2 * UDENS
GAMMA(IJ) = ODX2 * RDENS

```


***** GENPOI *****

173.	C	40			
174.	C	90	XP(IJ)	=	PRESSR(I,J)
175.	C	32			
176.		100	CORTINUE		
177.		200	CORTINUE		
178.	C	32			
179.		32	RETURN		
180.		32	END		

END ELT. ERRORS: NONE. TIME: 0.652 SEC. IMAGE COUNT: 180

@HDC,P ***** GENTS3 *****



***** HCCS IN *****

57. C CALLED AS 'CALL PSOLVE (PPARM, C, Y)';
 58. C PPARM IS A VECTOR OF PARAMETERS TO PSOLVE
 59. C C IS THE RIGHT HAND SIDE (WHICH MUST BE PRESERVED)
 60. C Y IS THE SOLUTION TO THE SYSTEM
 61. C
 62. C APARM - A VECTOR OF PARAMETERS TO THE SUBROUTINE ATIMES.
 63. C (USED TO OBIVIATE THE NEED FOR USE OF COMMON)
 64. C
 65. C PPARM - A VECTOR OF PARAMETERS TO THE SUBROUTINE PSOLVE.
 66. C PASSED DIRECTLY BY THIS SUBROUTINE TO PSOLVE.
 67. C
 68. C CNTL - IF > 1, PRINT SUMMARY OF ITERATIONS, OTHERWISE
 69. C SUPPRESS ALL OUTPUT
 70. C ON OUTPUT, 0 IS INDICATION OF SUCCESS. 1 INDICATES
 71. C FAILURE TO CONVERGE WITHIN ITERATION LIMIT.
 72. C
 73. C
 74. C
 75. C
 76. C
 77. C
 78. C
 79. C
 80. C
 81. C
 82. C
 83. C
 84. C
 85. C
 86. C
 87. C
 88. C
 89. C
 90. C
 91. C
 92. C
 93. C
 94. C
 95. C
 96. C
 97. C
 98. C
 99. C
 100. C
 101. C
 102. C
 103. C
 104. C
 105. C
 106. C
 107. C
 108. C
 109. C
 110. C
 111. C
 112. C
 113. C
 114. C

OUTPUT:
 X - THE APPROXIMATE SOLUTION TO AX = B. (SEE ABOVE)
 MAXITR - NUMBER OF ITERATIONS USED.

LALPHA, LBETA - THE COEFFICIENTS OF THE TRIDIAGONAL MATRIX WHICH WOULD HAVE BEEN PRODUCED BY THE LANCZOS ALGORITHM FOR FINDING EIGENVALUES. USEFUL FOR ESTIMATING THE CONVERGENCE RATES OF THE CONJUGATE GRADIENTS ALGORITHM, AND A CHEAP BYPRODUCT OF THE LATTER ALGORITHM.

WORKING PARAMETERS:
 (VECTORS FOR WORKING STORAGE WITHIN SUBROUTINE)
 R HOLDS APPROXIMATE (RECURSIVE) RESIDUAL FOR FINAL SOLUTION.
 P, AP TEMPORARIES NEEDED BY ALGORITHM.

 LOCAL VARIABLES:

INTEGER I, ITER
 REAL PTAP, RTPTR, NORMR, NORMB, AK, BK, XRTPIR, IP,
 EXTERNAL RELRES, AKHI
 LOGICAL IP
 INTEGER PRINTR
 DATA PRINTR / 6 /

 COMPUTE INITIAL RESIDUAL AND INITIAL SEARCH DIRECTION

FORCE LINEAR SYSTEM TO BE CONSISTENT BY REMOVING ANY COMPONENT OF THE NULLSPACE (ORTHOGONAL COMPLEMENT OF RANGE) FROM RIGHT HAND SIDE B. THEORETICALLY THE CONSISTENCY OF THE LINEAR SYSTEM THEN IMPLIES THAT THE VECTORS X AND R, WHICH ARE GENERATED BY THE RECURSION ARE ALL ORTHOGONAL TO THE NULLSPACE OF A. WE ENSURE THIS ORTHOGONALITY IN THE NUMERICAL



PROBLEM BY INCORPORATING A "REORTHOGONALIZATION". NOTE THAT THE VECTORS P ARE NOT NECESSARILY ORTHOGONAL TO THE NULLSPACE OF A.

```

115. C 07
116. C 07
117. C 07
118. C 07
119. C 07
120. C 07
121. C 09
122. C 06
123. C 07
124. C 06
125. C 06
126. C 07
127. C 06
128. C 06
129. C 07
130. C 07
131. C 07
132. C 06
133. C 09
134. C 06
135. C 06
136. C 06
137. C 07
138. C 10
139. C 06
140. C 06
141. C 06
142. C 06
143. C 06
144. C 06
145. C 06
146. C 06
147. C 06
148. C 06
149. C 09
150. C 06
151. C 06
152. C 06
153. C 06
154. C 06
155. C 06
156. C 06
157. C 06
158. C 06
159. C 07
160. C 07
161. C 07
162. C 06
163. C 06
164. C 06
165. C 06
166. C 06
167. C 06
168. C 06
169. C 06
170. C 06
171. C 06
172. C 06

```

CALL PROJECT (NM, B, NULLSP)

CALL ATIMES (APARM, X, AP)

NORMB = 0.0

NORMR = 0.9

DO 10 I = 1, NM

R(I) = B(I) - AP(I)

NORMR = NORMR + R(I)**2

NORMB = NORMB + B(I)**2

10 CONTINUE

RELRES = SQRT(NORMB/NORMB)

CALL PROJECT (NM, R, NULLSP)

CALL PSOLVE (PPARM, R, P)

RTPIR = IP (M1, P, R)

ITER = 6

PRINT = (CNTL .GE. 1)

IF (PRINT) WRITE (PRINTR, 6000) ITER, RELRES

M A I N I T E R A T I O N

ITER = ITER + 1

100 COMPUTE NEW SOLUTION GUESS AND NEW RESIDUAL

CALL ATIMES (APARM, P, AP)

PTAF = IP (NM, P, AP)

AK = RTPIR / PTAF

NORMR = 0.0

DO 110 I = 1, NM

X(I) = X(I) + AK * P(I)

R(I) = R(I) - AK * AP(I)

NORMR = NORMR + R(I)**2

110 CONTINUE

CALL PROJECT (NM, X, NULLSP)

CALL PROJECT (NM, R, NULLSP)

COMPUTE LANCZOS DIAGONAL ELEMENT

IF (ITER .GT. 1) GO TO 120

LALPHA(1) = 1.0 / AK

GO TO 130

LALPHA(ITER) = 1.0 / AK + BK / AICH

CONVERGENCE TEST


```

173. 06 RELRES = SQRT ( NORMR / NORMB)
174. 06 IF ( RELRES .LT. RTOL) .OR.
175. 06 ( ITER .EQ. MAXITER ) GO TO 200
176. 06
177. 06 COMPUTE NEW DIRECTION
178. 06
179. 06 AKM1 = AK
180. 09 CALL PSOLVE ( PFARM, R, AP)
181. 06 XRTPIR = IP ( NM, R, AP)
182. 06 BK = XRTPIR / RPIR
183. 06 IRTPIR = XRTPIR
184. 06
185. 06 DO 140 I = 1, NM
186. 06 P(I) = BK * P(I) + AP(I)
187. 06 CONTINUE
188. 07
189. 06
190. 06 COMPUTE LANCZOS OFF-DIAGONAL
191. 06
192. 06 LBETA(ITER+1) = SQRT(BK / AK
193. 10 IF (PRINT) WRITE (PRINT, 6050) ITER, RELRES, AK, BK,
194. 07 LALPHA(ITER), LBETA(ITER+1)
195. 06
196. 06 GO TO 100
197. 06
198. 06
199. 06
200. 10 IF (PRINT) WRITE (PRINT, 6100) ITER, RELRES, AK, LALPHA(ITER,
201. 10 ITER, RELRES, RTOL)
202. 07
203. 07 CNTL = 0
204. 07 IF (RELRES .GT. RTOL) CNTL = 1
205. 06 MAXITER = ITER
206. 06 RETURN
207. 06
208. 06
209. 06
210. 06
211. 07
212. 07
213. 07
214. 06
215. 06
216. 06
217. 07
218. 07
219. 06
220. 06
221. 06

```

END DLJ. ERRORS: NONE. TIME: 0.772 SEC. IMAGE COUNT: 221

***** HCCSOL *****

***** HCGSOL *****

@:ELT, L -PF1-.HCGSOL

ELT BR1 S7401C 07/11/79 08:48:46 (23)

SUBROUTINE HCGSOL

1 (M, N, NM, MMAX, FLIP, RESID,
 2 GENPO1, SCALE, BKSCAL, ATINES, PSOLVE,
 3 RHO, ELPDEF, X0, F, X,
 4 WRK1, WRK2, WRK3, NULLSP, APARM, ALPHA,
 5 BETA, GAMMA, LOWERB, UPPERB, LEFTB, RIGHTB,
 6 PPARM, DELTAX, DELTAY, MAXITR, RTOL, CNTL)

=====

SOLVE DISCRETE PRESSURE EQUATION
 GENERATE DISCRETE GENERALIZED POISSON OPERATOR (GENPO1)
 PERFORM DIAGONAL SCALING ON DISCRETE OPERATOR (SCALE)
 SOLVE SCALED LINEAR SYSTEM BY CONJUGATE GRADIENTS (HCGSIN)
 (USES FAST POISSON SOLVER OR OTHER SPLITTING (PSOLVE))
 BACK-TRANSFORM TO SOLUTION OF ORIGINAL PROBLEM (BKSCAL)

INTEGER M, N, NM, MMAX
 LOGICAL FLIP

SUBROUTINE PARAMETERS WHICH DESCRIBE THE TWO MATRICES WHICH APPEAR IN THE HYBRID CONJUGATE GRADIENT ALGORITHM

EXTERNAL GENPO1, SCALE, BKSCAL, ATINES, PSOLVE

ARRAYS ARE ALLOCATED IN A "DYNAMIC" SENSE IN THE DRIVER

ARRAYS OF DIMENSION MMAX BY N (PASSED FROM HYPERBOLIC PDE SOLVER)

REAL RHO(MMAX,N), ELPDEF(MMAX,N), X0(MMAX,N), RESID(MMAX,N)

ARRAYS OF DIMENSION M * N

REAL X(NM), F(NM), WRK1(NM), WRK2(NM), WRK3(NM), NULLSP(NM),
 ALPHA(NM), BETA(NM), GAMMA(NM)

... THE FOLLOWING ACTUAL PARAMETERS MAY OCCUPY THE SAME STORAGE

ELPDEF AND WRK1
 X0 AND WRK2
 RESID AND WRK3

ARRAY OF PARAMETERS FOR FIVE-DIAGONAL MATRIX MULTIPLY ROUTINE
 ACTUAL DIMENSION: 3NM + 2

REAL APARM(1)

ARRAYS OF DIMENSION M

REAL LOWERB(NM), UPPERB(NM)

ARRAYS OF DIMENSION N

19


```

57. REAL LEFTB(N), RIGHTB(N)
58.
59.
60.
61.
62.
63.
64.
65.
66.
67.
68.
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71.
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110.
111.
112.
113.
114.

ARRAY OF PARAMETERS FOR SOLVING SYSTEM WITH SPLITTING MATRIX
ACTUAL DIMENSION: 3L + 9K + 2(L+2)*(LOG(L)-1) + 12
WHERE (K,L) = (M,N) UNLESS N IS NOT ONE LESS THAN
A POWER OF TWO, OR N AND M ARE BOTH ONE LESS
THAN A POWER OF TWO, AND N IS LARGER
(IN WHICH CASE (K,L) = (N,M)).
THE INITIALIZING SUBROUTINE 'INTNEU' CHECKS THAT
SUFFICIENT SPACE IS AVAILABLE IN 'PPARM'.

REAL PPARAM(1)
SCALARS
REAL DELTAX, DELTAY, RTOL
INTEGER MAXITR, CNTL
=====
LOCAL VARIABLES ...
INTEGER LMXITR
REAL LALPHA(50), LBETA(50)
DIMENSION LMXITR
DATA LMXITR / 50 /

INTEGER I, J, I, K, L, ITER, ERROR
INTEGER TOTTIM, PRVTIM, STRTIM
REAL RVORM, FJRMX, ELPTIM, CONSST, DXDYSQ, COND, CONTRAT,
SFACTR

INTEGER PRINTR
DATA PRINTR / 6 /

=====
CALL CPUSUP (STRTIM)
PRINT = CNTL
IF (PRINT .GE. 3) WRITE (PRINTR, 3000) M, N
MAXITR = MINO(MAXITR, LMXITR)

GENERATE DISCRETIZATION OF PRESSURE EQUATION
CALL GENPO1 (M, N, NMAX, FLIP, THO, ELPDEF, X0, DELTAX, DELTAY,
LOWERB, UPPERB, LEFTB, RIGHTB,
ALPHA, BETA, GAMMA, F, X, WRK3)

```



```

115. 20 IF (FLIP) GO TO 10
116. 20 .. NATURAL ORDERING USED ..
117. 20
118. 20
119. 20 APARM(1) = M
120. 20 APARM(2) = N
121. 20 GO TO 20
122. 20
123. 20 .. TRANSPOSE OF NATURAL ORDERING USED ..
124. 20
125. 20 APARM(1) = N
126. 20 APARM(2) = M
127. 20
128. 20 K = APARM(1)
129. 20 L = APARM(2)
130. 20
131. 20 IF (PRINT .GE. 4) WRITE (PRINTR, 4000)
132. 20 IF (PRINT .GE. 4) CALL BTROUT (K, L, ALPHA, BETA, GAMMA)
133. 19
134. 20 FORCE LINEAR DIFFERENCE EQUATIONS TO BE CONSISTENT ...
135. 20 (EQUIVALENT TO REQUIRE THE RIGHT HAND SIDE TO HAVE MEAN ZERO)
136. 20 FORCE INITIAL PRESSURE GUESS TO HAVE MEAN PRESSURE ZERO
137. 20 (HENCE, IS ORTHOGONAL TO THE NULL SPACE OF THE LINEAR OPERATOR)
138. 20
139. 19 CALL PROJE (NM, F, CONSST)
140. 20 IF (PRINT .GE. 2) WRITE (PRINTR, 2000) CONSST
141. 19
142. 19 CALL PROJE (NM, X0, CONSST)
143. 20 IF (PRINT .GE. 2) WRITE (PRINTR, 2010) CONSST
144. 19
145. 20 COMPUTE RESIDUAL OF CONTINUOUS SOLUTION (OPTIONAL)
146. 19
147. 20 IF (PRINT .LT. 2) GO TO 210
148. 22 CALL ATIMES (APARM, X, WRK1)
149. 19 RNORM = 0.0
150. 19 NORFX = 0.0
151. 19 DO 100 I = 1, NM
152. 19 RNORM = RNORM + (F(I) - WRK1(I))**2
153. 19 NORFX = NORFX + X(I)**2
154. 19 WRK2(1) = 1.0
155. 19 CONTINUE
156. 19 RNORM = SQRT(RNORM/NORFX)
157. 20 WRITE (PRINTR, 2020) RNORM
158. 19
159. 20 IF (PRINT .GT. 2) WRITE (PRINTR, 3010)
160. 20 IF (PRINT .GT. 2) CALL VECOUT (K, L, F)
161. 19
162. 19 RNORM = 0.0
163. 22 CALL ATIMES (APARM, WRK2, WRK1)
164. 19 DO 200 I = 1, NM
165. 19 RNORM = RNORM + WRK1(I)**2
166. 20 CONTINUE
167. 19 RNORM = SQRT(RNORM/NM)
168. 20 WRITE (PRINTR, 2030) RNORM
169. 19
170. 19 SCALE DISCRETE OPERATOR
171. 19
172. 20 210 DXDYSQ = (DELTA/DELTAY) ** 2

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173. SFACTR = DDYDYSQ
174. IF (FLIP) SFACTR = 1.0 / SFACTR
175. CALL SCALE (K, L, ALPHA, BETA, GAMMA, SFACTR, NULLSP)
176.
177. IF (PRINT .LE. 3) GO TO 220
178. WRITE (PRINTR, 40:0)
179. CALL BTROUT (K, L, ALPHA, BETA, GAMMA)
180.
181.
182. SCALE RIGHT HAND SIDE OF EQUATION
183.
220 DO 300 I = 1, NM
    F(I) = F(I) / NULLSP(I)
    X(I) = X(I) * NULLSP(I)
300 CONTINUE
184.
185. CHECK ACCURACY OF SCALING AND SCALED NULLSPACE (OPTIONAL)
186.
187. IF (PRINT .LT. 2) GO TO 410
188. CALL ATIMES (APARM, NULLSP, WRK1)
189. RNORM = 0.0
190. NORMX = 0.0
191. DO 400 I = 1, NM
    RNORM = RNORM + WRK1(I)**2
    NORMX = NORMX + NULLSP(I)**2
400 CONTINUE
192. RNORM = SQRT(RNORM/NORMX)
193. NORMX = SQRT(NORMX)
194. WRITE (PRINTR, 2040) RNORM, NORMX
195.
196. CALL HYBRID CONJUGATE GRADIENTS ALGORITHM
197.
200. ITER = MAXITR
201.
202. CALL CPUSUP (TOTTIM)
203. ELPTIM = (TOTTIM - SIRTIM) / 1000.
204. PRVTIM = TOTTIM
205. IF (PRINT .GT. 1) WRITE (PRINTR, 2050) ELPTIM
206.
207. CALL HGSIN (NM, ITER, RTOL, X, F, WRK3, WRK1, WRK2, NULLSP,
    ATIMES, PSOLVE, APARM, PPARM, LALPHA, LBETA, CNTL)
208.
209. IF (PRINT .LE. 1) GO TO 420
210. CALL CPUSUP (TOTTIM)
211. ELPTIM = (TOTTIM - PRVTIM) / 1000.
212. WRITE (PRINTR, 2060) ELPTIM
213. PRVTIM = TOTTIM
214.
215. COMPUTE APPROXIMATE EIGENVALUES OF OPERATOR (OPTIONAL)
216.
217. CALL TQL1 (ITER, LALPHA, LBETA, ERROR)
218. IF (ERROR .NE. 0) WRITE (PRINTR, 2070) ERROR
219. COND = LALPHA(ITER) / LALPHA(1)
220. CONRAT = (COND - 1.0) / (COND + 1.0)
221. WRITE (PRINTR, 2080) COND, CONRAT
222.
223. SCALE THE PROBLEM BACK INTO THE ORIGINAL COORDINATES
224.
225.
226.
227.
228.
229.
230.

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```

231. 20 CALL BKSCAL (K, L, ALPHA, BETA, GAMMA, NULLSP)
232. 19 DO 500 I = 1, NM
233. 19 X(I) = X(I) / NULLSP(I)
234. 19 F(I) = F(I) * NULLSP(I)
235. 19 CONTINUE
236. 19
237. 19 C NORMALIZE SOLUTION TO MEAN PRESSURE ZERO (CONSISTENT)
238. 19 C
239. 19 C CALL PROJE (NM, X, CONSST)
240. 20 IF (PRINT .GE. 2) WRITE (PRINT, 2090) CONSST
241. 19 C
242. 19 C COMPUTE RESIDUALS FOR FINAL SOLUTION
243. 19 C AND REVERT TO TWO-DIMENSIONAL STORAGE
244. 19 C
245. 22 CALL ATIMES (APARM, X, WRK1)
246. 19 RNORM = 0.0
247. 19 NORMX = 0.0
248. 19 DO 600 I = 1, M
249. 19 DO 590 J = 1, N
250. 20 IF (.NOT. FLIP) IJ = (J-1)*M + I
251. 20 IF (FLIP) IJ = (I-1)*N + J
252. 20 RESID(I,J) = F(IJ) - WRK1(IJ)
253. 20 F(IJ) = RESID(I,J)
254. 19 RNORM = RNORM + RESID(I,J) ** 2
255. 19 X0(I,J) = X(IJ)
256. 19 NORMX = NORMX + X(IJ)**2
257. 19 CONTINUE
258. 19 C
259. 19 C
260. 19 C IF (PRINT .LT. 2) GO TO 700
261. 19 RNORM = SQRT(RNORM/NORMX)
262. 20 WRITE (PRINT, 1000) RNORM
263. 20 IF (PRINT .LT. 3) GO TO 610
264. 20 WRITE (PRINT, 3020)
265. 20 CALL VECOUT (K, L, F)
266. 20 WRITE (PRINT, 3030)
267. 20 CALL VECOUT (K, L, X)
268. 19 C
269. 20 C 610 CALL CPUSUP (TOTTIM)
270. 19 ELPTIM = (TOTTIM-STARTIM) / 1000.
271. 20 WRITE (PRINT, 2100) ELPTIM
272. 19 C
273. 20 C 700 RETURN
274. 19 C
275. 19 C
276. 19 C
277. 19 C
278. 20 C ARRANGED BY PRINT LEVEL ... FIRST DIGIT OF FORMAT NUMBER
279. 20 C INDICATES THE LEVEL AT WHICH IT BECOMES USED.
280. 20 C
281. 20 C 1000 FORMAT (' RELATIVE NORM OF RESIDUAL VECTOR A X - B :', 1PE10.2)
282. 20 C
283. 20 C 2000 FORMAT (' MEAN OF RIGHT HAND SIDE (REMOVED FOR CONSISTENCY) :',
284. 20 C 1 1PE12.3)
285. 20 C
286. 20 C 3010 FORMAT (' MEAN OF INITIAL PRESSURE APPROXIMATION :',
287. 20 C 1 1PE12.3)
288. 20 C

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289. C 2020 FORMAT (' RESIDUAL FOR UNSCALED CONTINUOUS PROB:', IPE12.3)
290. C 2030 FORMAT (' NORM OF OPERATOR APPLIED TO "E"', IPE12.3)
291. C 2040 FORMAT (' SCALED OPERATOR ON SCALED NULLSPACE', IPE12.3 /
292. C 1, ' NORM OF SCALED NULLSPACE', E12.3)
293. C 2050 FORMAT (' OSET UP TIME (SECONDS)', F13.3)
294. C 2060 FORMAT (' OTIME IN CONJUGATE GRADIENTS(SECONDS)', F13.3)
295. C 2070 FORMAT (' TOLL FAILURE ... SOME EIGENVALUES NOT COMPUTED' /
296. C 1, ' ERROR CODE RETURNED:', I5)
297. C 2080 FORMAT (' CONDITION NUMBER OF ITERATION MATRIX:', IPE12.3 /
298. C 1, ' ASYMPTOTIC CONVERGENCE RATE:', E12.3)
299. C 2090 FORMAT (' MEAN PRESSURE OF SOLUTION (REMOVED):', IPE12.3)
300. C 2100 FORMAT (' TOTAL TIME FOR RUN:', F13.3)
301. C 3000 FORMAT (' OPRESSURE EQUATION TEST' /
302. C 2, ' OGRID SIZE:', /
303. C 3, ' M:', I3 /
304. C 4, ' N:', I3)
305. C 3010 FORMAT (' ORIGHT HAND SIDE:')
306. C 3020 FORMAT (' ORESIDUAL VECTOR')
307. C 3030 FORMAT (' O' / ' OFINAL SOLUTION')
308. C 4000 FORMAT (' O' / ' UNSCALED DISCRETE OPERATOR ... ' /)
309. C 4010 FORMAT (' O' / ' OUNSCALED DISCRETE OPERATOR ... ' /)
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END ELT. ERRORS: NONE. TIME: 1.128 SEC. IMAGE COUNT: 326

CHDG, P ***** INTDIR *****

***** INTNEU *****

0:ELT,L -PP1-.INTNEU 08:48:54 (21) SUBROUTINE INTNEU (PPARM, LPPARM, M, N, DXDYSQ)

1. 18 C
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INITIALIZE PARAMETER ARRAY FOR SOLVING THE DISCRETE LAPLACIAN ON A RECTANGLE WITH NEUMANN BOUNDARY CONDITIONS, USING THE NCAR ROUTINE 'BLKTR1' BY PAUL SCHWARZ/TRAUBER.

=====
THIS ROUTINE MUST BE CALLED ONCE, BEFORE THE ELLIPTIC PDE SOLVER 'HCSOL' IS USED. IT DOES NOT NEED TO BE REINITIALIZED AGAIN.
=====

INTEGER LPPARM, M, N
REAL FPARM(LPPARM)
REAL DXDYSQ

PPARM IS AN ARRAY CONTAINING ALL THE PARAMETERS NEEDED BY THE SUBROUTINE NEUSOL. THE INDIVIDUAL ELEMENTS ARE NAMED AND COMMENTED IN THE CODE SETTING UP THE ARRAY VALUES. PPARM MUST BE A VECTOR OF LENGTH AT LEAST
 $3N + 9M + 2(N+2)(\text{LOG}(N)-1) + 12$
OR
 $3M + 9N + 2(M+2)(\text{LOG}(M)-1) + 12$
WHERE THE LOGARITHM IS BASE 2, AND N (ND) MUST BE ONE LESS THAN A POWER OF 2.

LPPARM IS THE DECLARED LENGTH OF PPARM (THE INTEGER CONSTANT IN THE REAL OR DIMENSION STATEMENT.) THIS SUBROUTINE WILL CALCULATE THE ACTUAL STORAGE NEEDED AND WILL STOP IF THERE IS NOT ENOUGH.

M IS THE NUMBER OF GRID POINTS IN THE X DIRECTION,
N IS THE NUMBER OF GRID POINTS IN THE Y DIRECTION.

DXDYSQ IS THE SQUARE OF THE RATIO DELTA X / DELTA Y.

INTEGER I, SUBAL, SUBAK, SUBBL, SUBBK, SUBCL, SUBCK, SUBWK,
LOGICAL INIT, NEED, K, L
REAL FLIP, SFACTR

INTEGER PRINTR
DATA PRINTR / 6 /

DETERMINE WHICH ORDER IS MORE APPROPRIATE FOR THE FAST SOLVER, AT LEAST ONE OF THE LIMITS M AND N MUST BE ONE LESS THAN A POWER OF TWO. L WILL BE THE SMALLER OF M AND N WHICH MEETS THIS CONDITION.

FLIP = ((N .NE. 2**((LOG2(N+1)) - 1)) .OR.
1 ((N .EQ. 2**((LOG2(N+1)) - 1)) .AND.


```

57. 18 (M .EQ. 2***(LOG2(M+1)) - 1) .AND.
58. 16 (M .LT. N) )
59. 18
60. 18
61. 18
62. 18 IF (FLIP) GO TO 10
63. 18
64. 13 .. NATURAL ORDERING USED ..
65. 18
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106. 18
107. 12
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111. 18
112. 18
113. 16
114. 18

```

2 (M .EQ. 2***(LOG2(M+1)) - 1) .AND.
 3 (M .LT. N))
 IF (FLIP) GO TO 10
 .. NATURAL ORDERING USED ..
 K = M
 L = N
 GO TO 20
 .. TRANSPOSED ORDERING USED ..
 10 K = N
 L = M
 .. CHECK STORAGE ALLOCATION ..
 NEED = 3*L + 9*K + 2*(L+2)*(LOG2(L+1)-1) + 12
 20 IF (LPPARM .LT. NEED) GO TO 1000
 INITIALIZATION FLAG:
 PPARM(1) = 0.0
 X DIMENSION
 PPARM(2) = K
 Y DIMENSION
 PPARM(3) = L
 LOCATION OF AK
 SUBAK = 11
 PPARM(4) = SUBAK
 LOCATION OF AL
 SUBAL = SUBAK + K
 PPARM(5) = SUBAL
 LOCATION OF BK
 SUBBK = SUBAL + L
 PPARM(6) = SUBBK
 LOCATION OF BL
 SUBBL = SUBBK + K
 PPARM(7) = SUBBL
 LOCATION OF CK
 SUBCK = SUBBL + L
 PPARM(8) = SUBCK
 LOCATION OF CL
 SUBCL = SUBCK + K
 PPARM(9) = SUBCL
 LOCATION OF WORKING STORAGE
 PPARM(10) = SUBCL + L

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172. C

INITIALIZE CONSTANTS FOR DIAGONALS

SFACTR = DXDYSD
IF (FLIP) SFACTR = 1.0 / SFACTR

PPARM(SUBAKO) = 0
PPARM(SUBAL) = 0
PPARM(SUBBK) = -1.0
PPARM(SUBEL) = -SFACTR
PPARM(SUBCK) = 1.0
PPARM(SUBCL) = SFACTR

DO 100 I = 3, K
SUBAK = SUBAK + 1
SUBBK = SUBBK + 1
SUBCK = SUBCK + 1
PPARM(SUBAKO) = 1.0
PPARM(SUBBK) = -2.0
PPARM(SUBCK) = 1.0

100 CONTINUE

PPARM(SUBAK+1) = 1.0
PPARM(SUBBK+1) = -1.0
PPARM(SUBCK+1) = 0.0

DO 200 I = 3, L
SUBAL = SUBAL + 1
SUBBL = SUBBL + 1
SUBCL = SUBCL + 1
PPARM(SUBAL) = SFACTR
PPARM(SUBEL) = -2.0 * SFACTR
PPARM(SUBCL) = SFACTR

200 CONTINUE

PPARM(SUBAL+1) = SFACTR
PPARM(SUBBL+1) = -SFACTR
PPARM(SUBCL+1) = 0.0

INITIALIZE FAST POISSON SOLVER

SUBAK = PPARM(4)
SUBAL = PPARM(5)
SUBBK = PPARM(6)
SUBEL = PPARM(7)
SUBCK = PPARM(8)
SUBCL = PPARM(9)
SUBMK = PPARM(10)
INIT = 0

CALL BLKTR: (INIT, 1, L, PPARM(SUBAL), PPARM(SUBBL),
1 PPARM(SUBCL),
2 1, K, PPARM(SUBAK), PPARM(SUBBK),
3 PPARM(SUBCK), K,
4 PINVX, ERROR, PPARM(SUBMK))

IF (ERROR .NE. 0) GO TO 2000


```

173.          PPARAM(1) = 1.0
174.          RETURN
175.          C
176.          C
177.          C
178.          C
179.          C
180.          C
181.          C
182.          C
183.          C
184.          C
185.          C
186.          C
187.          C
188.          C
189.          C
190.          C
191.          C
192.          C
193.          C
194.          C
195.          C

1000 WRITE (PRINTR, 6000) NEED, LPPARM
STOP
6000 FORMAT ('0***** NOT ENOUGH PARAMETER STORAGE FOR BLKTRI' /
, ' PPARAM MUST BE AT LEAST', 14, ' LONG' /
, ' LPPARM WAS ONLY', 110 /
, ' *****')
C
2000 WRITE (PRINTR, 6100) ERROR
STOP
C
6100 FORMAT ('0***** FATAL ERROR IN INITIALIZING POISSON SOLVER' /
, ' ERROR RETURN FLAG:', 13 /
, ' CHECK WRITEUP OF NCAR ROUTINE BLKTRI' /
, ' FOR INTERPRETATION OF FLAG' /
, ' *****')
END

```

END ELT. ERRORS: NONE. TIME: 0.703 SEC. IMAGE COUNT: 195

@HDC,P ***** IP *****

***** IP *****

@:ELT,L -P71-.IP

ELT @RI S7401C 07/11/79 08:48:56 (P)

REAL FUNCTION IP (N, X, Y)

- 1. 00 C
- 2. 00 C
- 3. 00 C
- 4. 00 C
- 5. 00 C
- 6. 00 C
- 7. 00 C
- 8. 00 C
- 9. 00 C
- 10. 00 C
- 11. 00 C
- 12. 00 C
- 13. 00 C
- 14. 00 C

COMPUTE THE INNERPRODUCT OF TWO REAL VECTORS

INTEGER N, I
REAL X(N), Y(N)

IP = 0.0

DO 10 I = 1, N

IP = IP + X(I) * Y(I)

10 CONTINUE

RETURN

END

END ELT. ERRORS: NONE. TIME: 0.109 SEC. IMAGE COUNT: 14

@HDC,P ***** LISTPDE *****

@:ELT,L -PF1-.NEUBAK
ELT BR1 S74Q1C 07/11/79

03:49:06 (3)
SUBROUTINE NEUBAK (M, N, ALPHA, BETA, GAMMA, NULLSP)

```

1. 03 C
2. 00 C
3. 00 C
4. 02 C
5. 02 C
6. 02 C
7. 02 C
8. 02 C
9. 03 C
10. 00 C
11. 00 C
12. 00 C
13. 00 C
14. 03 C
15. 00 C
16. 00 C
17. 00 C
18. 00 C
19. 00 C
20. 00 C
21. 02 C
22. 00 C
23. 00 C
24. 00 C
25. 00 C
26. 00 C
27. 00 C
28. 00 C
29. 03 C
30. 02 C
31. 00 C
32. 00 C
33. 00 C
34. 00 C
35. 00 C
36. 00 C
37. 00 C
38. 00 C
39. 00 C
40. 00 C

```

REVERSE THE SCALING CARRIED OUT BY 'NEUSCL'. THERE WE
 COMPUTED A SCALING TO REDUCE THE MAIN DIAGONAL OF A SPARSE
 MATRIX TO BE THAT OF THE DISCRETE LAPLACIAN-NEUMANN PROBLEM.
 WE WANT NOW:
 A = D***.5 B D***.5
 WHERE THE VECTOR NULLSP HOLDS THE VALUES OF D***.5

INTEGER M, N
 REAL ALPHA(1), BETA(1), GAMMA(1), NULLSP(1)

THE VECTOR PARAMETERS ARE ONE-DIMENSIONAL ORDERINGS OF THE
 TWO DIMENSIONAL GRID PROBLEM. THE ACTUAL DIMENSION OF
 THE VECTORS IS 'N' * 'M'.

INTEGER K, XI, YJ
 REAL SCALE

```

K = 1
DO 200 YJ = 1, N
DO 100 XI = 1, M
SCALE = NULLSP(K)
ALPHA(K) = ALPHA(K) * (SCALE**2)
BETA(K) = BETA(K) * SCALE
GAMMA(K) = GAMMA(K) * SCALE
IF ( YJ .GT. 1 ) GAMMA(K-M) = GAMMA(K-M) * SCALE
IF ( XI .GT. 1 ) BETA(K-1) = BETA(K-1) * SCALE
K = K + 1
100 CONTINUE
200 CONTINUE
RETURN
END

```

END ELT. ERRORS: NONE. TIME: 0.214 SEC. IMAGE COUNT: 40

@HDC,P ***** NEUSCL *****

6:ELT,L -PF1-.NEUSCL
 ELT 8R1 S74Q1C 07/11/79 06:49:07 (7)
 SUBROUTINE NEUSCL (M, N, ALPHA, BETA, GAMMA, DXDYSQ, NULLSP)

```

1. 07 C
2. 04 C
3. 04 C
4. 04 C
5. 04 C
6. 04 C
7. 04 C
8. 04 C
9. 04 C
10. 04 C
11. 04 C
12. 04 C
13. 07 C
14. 07 C
15. 07 C
16. 07 C
17. 04 C
18. 04 C
19. 04 C
20. 04 C
21. 07 C
22. 04 C
23. 04 C
24. 04 C
25. 04 C
26. 04 C
27. 04 C
28. 04 C
29. 04 C
30. 04 C
31. 04 C
32. 04 C
33. 04 C
34. 04 C
35. 04 C
36. 04 C
37. 04 C
38. 04 C
39. 04 C
40. 05 C
41. 04 C
42. 06 C
43. 04 C
44. 04 C
45. 07 C
46. 04 C
47. 04 C
48. 04 C
49. 04 C
50. 04 C
51. 04 C
52. 04 C
53. 04 C
54. 04 C
55. 04 C
56. 04 C

```

MULTIPLY THE SPARSE SYMMETRIC MATRIX WITH FIVE NON-ZERO DIAGONALS (ALPHA, BETA, GAMMA) BY A SCALING MATRIX D CHOSEN SO THE MAIN DIAGONAL REDUCES TO THAT OF THE DISCRETE LAPLACIAN WITH NEUMANN BOUNDARY CONDITIONS. IF A IS THE ORIGINAL MATRIX, WE COMPUTE

B = D**(-.5) A D**(-.5)

WE RETURN THE DIAGONALS OF B IN THE PLACE OF THE DIAGONALS OF A. THE VECTOR NULLSP RETURNS THE VALUES OF THE NON-ZERO ELEMENTS OF D**-.5, WHICH GIVES THE NULLSPACE OF THE SCALED OPERATOR (ASSUMING THE ORIGINAL OPERATOR WAS DERIVED FROM A FIVE-POINT FORMULA, WITH NULLSPACE (1,1,1,....,1))

```

-----
INTEGER M, N
REAL ALPHA(1), BETA(1), GAMMA(1), NULLSP(1), DXDYSQ

```

THE VECTOR PARAMETERS ARE ONE-DIMENSIONAL ORDERINGS OF THE TWO DIMENSIONAL GRID PROBLEM. THE ACTUAL DIMENSION OF THE VECTORS IS 'N' * 'M'.

```

INTEGER K, XI, YJ
REAL FACTOR, SCALE

```

TWO DIMENSIONAL ORDERING USED IMPLICITLY TO FIND AND ADJUST THE ELEMENTS CORRESPONDING TO BOUNDARY POINTS IN THE GRID

```

K = 1
DO 200 YJ = 1, N
  DO 100 XI = 1, M
    FACTOR = -2.0 * (1.0 + DXDYSQ)
    IF ((XI .EQ. 1) .OR. (XI .EQ. M)) FACTOR = FACTOR + 1.0
    IF ((YJ .EQ. 1) .OR. (YJ .EQ. N)) FACTOR = FACTOR + DXDYSQ

```

```

SCALE = SQRT (FACTOR / ALPHA(K))
NULLSP(K) = 1.0 / SCALE
ALPHA(K) = FACTOR
BETA(K) = BETA(K) * SCALE
GAMMA(K) = GAMMA(K) * SCALE
IF ( YJ .GT. 1 ) GAMMA(K-M) = GAMMA(K-M) * SCALE
IF ( XI .GT. 1 ) BETA(K-1) = BETA(K-1) * SCALE
K = K + 1

```

```

100 CONTINUE
200 CONTINUE
RETURN
END

```


***** NEUSOL *****

END ELT. ERRORS: NONE. TIME: 0.269 SEC. IMAGE COUNT: 56

@HDC,P ***** NEUSOL *****

DATE 071179

PAGE 2

***** NEUSOL *****

Q:ELT,L -PF1-.NEUSOL
ELT BR1 S74QIC 07/11/79

03:49:10 (32)

SUBROUTINE NEUSOL (PPARM, X, PINVX)

1 C
2 C
3 C
4 C
5 C
6 C
7 C
8 C
9 C
10 C
11 C
12 C
13 C
14 C
15 C
16 C
17 C
18 C
19 C
20 C
21 C
22 C
23 C
24 C
25 C
26 C
27 C
28 C
29 C
30 C
31 C
32 C
33 C
34 C
35 C
36 C
37 C
38 C
39 C
40 C
41 C
42 C
43 C
44 C
45 C
46 C
47 C
48 C
49 C
50 C
51 C
52 C
53 C
54 C
55 C
56 C

SOLVE P Y = X
WHERE P IS THE DISCRETE NEUMANN LAPLACIAN

REAL PPARM(1), X(1), PINVX(1)

PPARM IS A LIST OF PARAMETERS FOR THIS ROUTINE.
THIS ROUTINE UNPACKS THEM TO PASS THEM TO BLKTRI

INTEGER INIT, N, M, SUBAM, SUBAN, SUBBM, SUBBN, SUBCM, SUBCN,
SUBCN, SUBWRK, NM, I, ERROR

REAL ADUMDY

INTEGER PRINTR
DATA PRINTR / 6 /

UNPACKING ...

INIT = IFIX(PPARM(1))
M = IFIX (PPARM(2))
N = IFIX (PPARM(3))

SUBAM = IFIX (PPARM(4))
SUBAN = IFIX (PPARM(5))
SUBBM = IFIX (PPARM(6))
SUBBN = IFIX (PPARM(7))
SUBCM = IFIX (PPARM(8))
SUBCN = IFIX (PPARM(9))
SUBWRK = IFIX (PPARM(10))

... COPY RIGHT HAND SIDE VECTOR X TO PROTECT IT ...

NM = N * M
DO 10 I = 1, NM
PINVX(I) = X(I)
10 CONTINUE

FORCE LINEAR SYSTEM TO BE CONSISTENT

CALL PROJCE (NM, PINVX, ADUMDY)

INVOKE NCAR ROUTINE ...

IF (INIT .NE. 1) GO TO 200


```

57. 31 CALL BLKTRI ( INIT, I, N, PPARM(SUBAN), PPARM(SUBBN),
58. 31 1 PPARM(SUBCN), PPARM(SUBBN), PPARM(SUBBN),
59. 31 1, M, PPARM(SUBAN), PPARM(SUBBN), PPARM(SUBBN),
60. 31 3 PPARM(SUBCN), M,
61. 31 PPARM(SUBCN), M,
62. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
63. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
64. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
65. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
66. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
67. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
68. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
69. 32 PPARM(SUBCN), M, PPARM(SUBCN), M,
70. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
71. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
72. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
73. 32 PPARM(SUBCN), M, PPARM(SUBCN), M,
74. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
75. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
76. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
77. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
78. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
79. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,
80. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
81. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
82. 29 PPARM(SUBCN), M, PPARM(SUBCN), M,
83. 28 PPARM(SUBCN), M, PPARM(SUBCN), M,

```

END ELT. ERRORS: NONE. TIME: 0.331 SEC. IMAGE COUNT: 33

@HDC,P ***** PROJCE *****


```

9:ELT,L -PF1-.PROJECT
ELT 3R1 S74Q1C 07/11/79 08:49:10 (7)
1. 04 C SUBROUTINE PROJECT (N, X, ORTHOG)
2. 03 C
3. 03 C PROJECT THE N VECTOR X ONTO THE ORTHOGONAL
4. 04 C COMPLEMENT OF THE VECTOR E.
5. 03 C
6. 03 C
7. 03 C
8. 03 C
9. 04 C
10. 03 C INTEGER N
11. 03 C REAL X(N), ORTHOG
12. 07 C INTEGER I
13. 03 C REAL IPXL
14. 04 C IPXE = 0.0
15. 03 C
16. 03 C DO 100 I = 1, N
17. 04 C IPXE = IPXE + X(I)
18. 03 C
19. 03 C 100 CONTINUE
20. 04 C ORTHOG = IPXE / N
21. 03 C
22. 03 C DO 200 I = 1, N
23. 04 C X(I) = X(I) - ORTHOG
24. 03 C 200 CONTINUE
25. 03 C
26. 03 C RETURN
27. 03 C
=====

```

END ELT. ERRORS: NONE. TIME: 0.151 SEC. IMAGE COUNT: 27

@MDG,P ***** PROJECT *****

ELT, L -PF1-.PROJECT
 ELT BR1 S74Q1C 07/11/79

08:49:12 (8)
 SUBROUTINE PROJECT (N, X, P)

PROJECT THE N VECTOR X ONTO THE ORTHOGONAL
 COMPLEMENT OF THE VECTOR P.

=====

INTEGER N
 REAL X(N), P(N)

INTEGER I
 REAL NORMP, IPXP, ORTHOG

NORMP = 0.0
 IPXP = 0.0

DO 100 I = 1, N
 NORMP = NORMP + P(I)**2
 IPXP = IPXP + X(I) * P(I)

100 CONTINUE

ORTHOG = IPXP / NORMP

DO 200 I = 1, N
 X(I) = X(I) - ORTHOG*P(I)

200 CONTINUE

RETURN
 END

END ELT. ERRORS: NONE. TIME: 0.157 SEC. IMAGE COUNT: 29

@HDC,P ***** VECOUT *****

@:ELT.L -PF1-.VECOUT

ELT 8R1 S7401C 07/11/79 00:49:13 (7)

SUBROUTINE VECOUT (M, N, VEC)

OUTPUT A BLOCK VECTOR IN SEMI-READABLE FORM

M : X DIMENSION GRID POINTS
N : Y DIMENSION GRID POINTS

INTEGER M, N
REAL VEC(M,N)

INTEGER MAXWID, NPAGES, P, START, FINISH
DATA MAXWID / 8 /

INTEGER PRINTR
DATA PRINTR / 6 /

=====

NPAGES = M / MAXWID .LT. M NPAGES = NPAGES + 1
IF (NPAGES * MAXWID .LT. M) NPAGES = NPAGES + 1

DO 100 P = 1, NPAGES
START = (P-1) * MAXWID + 1
FINISH = MIN1(P*MAXWID, M)
DO 50 J = 1, N
WRITE (PRINTR, 6000) (VEC(I,J), I = START, FINISH)

50 CONTINUE
WRITE (PRINTR, 6100)
100 CONTINUE

RETURN

=====

6000 FORMAT (' ', 1P9E9.2)

6100 FORMAT (' ')

END

END ELT. ERRORS: NONE. TIME: 0.161 SEC. IMAGE COUNT: 40

@HDC,N

@OFF

U.S. DEPT. OF COMM. BIBLIOGRAPHIC DATA SHEET <i>(See instructions)</i>	1. PUBLICATION OR REPORT NO. NBSIR 80-2056 (NBS)	2. Performing Organ. Report No.	3. Publication Date April 16, 1980
4. TITLE AND SUBTITLE The Numerical Solution of a Nonseparable Elliptic Partial Differential Equation by Preconditioned Conjugate Gradients			
5. AUTHOR(S) John Gregg Lewis and Ronald G. Rehm			
6. PERFORMING ORGANIZATION <i>(If joint or other than NBS, see instructions)</i> NATIONAL BUREAU OF STANDARDS DEPARTMENT OF COMMERCE WASHINGTON, D.C. 20234		7. Contract/Grant No.	8. Type of Report & Period Covered Final
9. SPONSORING ORGANIZATION NAME AND COMPLETE ADDRESS <i>(Street, City, State, ZIP)</i> Center for Applied Mathematics National Engineering Laboratory National Bureau of Standards			
10. SUPPLEMENTARY NOTES <input type="checkbox"/> Document describes a computer program; SF-185, FIPS Software Summary, is attached.			
11. ABSTRACT <i>(A 200-word or less factual summary of most significant information. If document includes a significant bibliography or literature survey, mention it here)</i> <p>In this report the combination of an iterative technique, the conjugate gradient algorithm, with a fast direct method, cyclic reduction, is used to solve the linear algebraic equations resulting from discretization of a nonseparable elliptic partial differential equation. An expository discussion of the conjugate gradient and preconditioned conjugate gradient algorithms and of their use in the solution of partial differential equations is presented. New results extending the use of the preconditioned conjugate gradients technique to singular linear equations which arise from discretized elliptic equations with Neumann boundary conditions are also given. The algorithms are applied to solve a specific elliptic equation which arises in the study of buoyant convection produced by a room fire. A code was developed to implement the algorithms for this application. Numerical results obtained through testing and use of the code are discussed.</p>			
12. KEY WORDS <i>(Six to twelve entries; alphabetical order; capitalize only proper names; and separate key words by semicolons)</i> conjugate gradient algorithm; elliptic partial differential equations; iterative methods for linear algebraic equations; Neumann boundary conditions; space matrices.			
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