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Bound for the *P*-Condition Number of Matrices With Positive Roots

Philip J. Davis, Emilie V. Haynsworth, and Marvin Marcus

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Let a matrix A have positive roots $0 < \lambda_1 \leq \lambda_2 \leq \ldots \leq \lambda_n$. Upper and lower bounds for the *P*-condition number of A, $P = \lambda_n / \lambda_1$, are given in terms of det A and one other symmetric function of the roots.

1. Introduction

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ppose
$$A = (a_{ij})$$
 is a nonsingular $n \times n$ matrix $c_k = a_{ij}$ roots, $\lambda_i (i=1, \ldots, n)$, ordered so that

$$0 < |\lambda_1| \le |\lambda_2| \le \ldots \le |\lambda_n|.$$

For a wide class of matrices A, the ratio $|\lambda_n|/|\lambda_1|$ gives a rough measure of the probable accuracy of the computation of the inverse of A, or the solution of the system of equations for which A is the matrix of coefficients.

This measure was evaluated in some detail by von Neumann and Goldstine $[5]^1$ and has been called by J. Todd [2, 3, 4] the *P*- condition number of the matrix, i.e.,

$$P = \frac{|\lambda_n|}{|\lambda_1|}$$

In general the accuracy of the results is in proportion to the reciprocal of P.

In this paper we shall show that, if the roots of A are all positive with

$$0 < \lambda_1 \le \lambda_2 \le \dots \le \lambda_n \tag{1}$$

and the determinant of A, det A, is known, together with one other symmetric function of the roots, upper and lower bounds for P can be given in terms of these two known quantities.

More precisely, if the characteristic polynomial for the matrix A is

$$p(x) = x^{n} - c_{1}x^{n-1} + c_{2}x^{n-2}$$

-...+(-1)ⁿc_n, c_i \ge 0 (i=1, ..., n) (2)

and we know c_k and $c_n = (\det A)$, there exist positive upper and lower bounds for the ratio

$$P = \frac{\lambda_n}{\lambda_1},\tag{3}$$

in terms of c_k and c_n .

To simplify the notation in the statement and proof of the theorem we use the following device.

$$c_k = c_k(\lambda_1, \ldots, \lambda_n) = \binom{n}{k} s_k(\lambda_1, \ldots, \lambda_n) = \binom{n}{k} s_k; \quad (4)$$

and divide each root of p(x) by $(s_k)^{1/k}$. (This is equivalent to dividing each element of the matrix by $(s_k)^{1/k}$.) We will call the new matrix "normalized with respect to k" or simply the normalized matrix. Since the condition number, P, is the ratio of two roots, P will not be affected by such a transformation. Let D_k be the determinant of the normalized matrix. Then

$$D_{k} = \frac{s_{n}}{(s_{k})^{n/k}} = \frac{\binom{n}{k}^{n/k} c_{n}}{c_{k}^{n/k}}.$$
(5)

From Hardy-Littlewood-Pólya [1] the numbers s_k as defined in (4) satisfy the inequalities

$$s_1 \ge s_2^{1/2} \ge s_3^{1/3} \ge \ldots \ge s_n^{1/n}$$
 (6)

(where equality holds only if all the λ_i are equal) hence, by (5), $D_k \leq 1$.

2. Statement and Proof of Theorem

THEOREM: If we have a matrix $A=(a_{ij})$ with characteristic polynomial (2) and if the constant term and the kth coefficient are known, the following bounds hold for P,

$$\frac{1}{D_k^{1/(n-1)}} \le P \le \frac{1 + \sqrt{1 - D_k^{n-1}}}{1 - \sqrt{1 - D_k^{n-1}}}, \qquad (k = 1, \dots, n-1) \quad (7)$$

where D is defined by (4) and (5).

If k=1 (i.e., the trace of the matrix is given) we can improve the upper bound:

$$P \leq \frac{1 + \sqrt{1 - D_1}}{1 - \sqrt{1 - D_1}}; \qquad D_1 = \frac{n^n c_n}{c_1^n}. \tag{8}$$

PROOF: Part 1—Upper bound. We prove (8) first, as the method used for this can be applied in the proof of (7).

Since $P = \lambda_n / \lambda_1$, eq (8) is equivalent to

$$D_1 \le \frac{4\lambda_1\lambda_n}{(\lambda_1 + \lambda_n)^2} \tag{9}$$

¹ Figures in brackets indicate the literature references at the end of this paper.

so from (5), we must show that

$$\frac{(\lambda_1 + \lambda_n)^2}{4\lambda_1\lambda_n} \le \left(\frac{s_1}{n\sqrt{s_n}}\right)^n \cdot$$
(10)

But the right side of (10) is $(A_n/G_n)^n$ where A_n is the arithmetic mean, and G_n the geometric mean, of $\lambda_1, \ldots, \lambda_n$.

Thus, we can prove (10) by induction if we prove the following:

LEMMA:

$$\left(\frac{A_n}{\overline{G}_n}\right)^n \le \left(\frac{A_{n+1}}{\overline{G}_{n+1}}\right)^{n+1},\tag{11}$$

where A_{n+1} and G_{n+1} are the arithmetic and geometric means of $\lambda_1, \ldots, \lambda_n, \lambda_{n+1}$. (The λ_i are not assumed to be ordered with respect to size for this lemma.)

PROOF OF LEMMA: Simplifying (11) we see we must prove

$$\lambda_{n+1}A_n^n \le A_{n+1}^{n+1}. \tag{12}$$

But (12) follows from the arithmetic-geometric mean inequality (cf., Hardy, Littlewood, Polya [1]),

$$a_1a_2 \ldots a_{n+1} \leq \left(\frac{a_1 + a_2 + \ldots + a_{n+1}}{n+1}\right),^{n+1}$$

if we let $a_i = A_n$, $(i=1, \ldots, n)$ and $a_{n+1} = \lambda_{n+1}$. Thus (11) holds and (10) follows immediately

using finite induction on n.

To prove (7) we note that, by the inequalities (6),

$$D_k^{n-1} \le D_{n-1}^{n-1} = \frac{s_n^{n-1}}{s_{n-1}^n},$$

so it suffices to show that

$$\frac{(\lambda_1+\lambda_n)^2}{4\lambda_1\lambda_n} \leq \frac{1}{D_{n-1}^{n-1}} = \left(\frac{s_{n-1}}{s_n}\right)^n \cdot s_n.$$
(13)

Now let $\mu_1 = \lambda_n^{-1}, \ldots, \mu_n = \lambda_1^{-1}$ and we note that $0 < \mu_1 \leq \ldots \leq \mu_n$ and (13) becomes

$$\frac{(\mu_1 + \mu_n)^2}{4\mu_1\mu_n} \leq \left(\frac{s_1(\mu_1 \dots \mu_n)}{s_n^{1/n}(\mu_1 \dots \mu_n)}\right)^n \cdot$$
(14)

But (14) is the same as (10) with the μ 's substituted for the λ 's, hence (13) holds and the general upper bound (7) is true for all values of k.

PROOF: Part 2-Lower Bound. From the inequalities (6),

$$D_k = \frac{s_n}{s_k^{n/k}} \ge \frac{s_n}{s_1^n},$$

where, by the definition (4)

Hence

$$\frac{1}{D_k} \leq \frac{s_1^n}{s_n} \leq \frac{\lambda_n^n}{\lambda_1^{n-1}\lambda_n} = \left(\frac{\lambda_n}{\lambda_1}\right)^{n-1}$$

 $s_1 \leq \lambda_n, s_n \geq \lambda_1^{n-1} \lambda_n.$

which proves the lower bound in (7).

3. Some Remarks

In reference [6] T. Kato gives an upper bound for the *P*-condition number which in our notation is

$$P < rac{4}{D_1}$$

This is larger than the bound given in (8).

We derive several consequences from (7), (8). Suppose that we have a matrix with positive roots, which is normalized so that $c_k = \binom{n}{k}$, i.e., $s_k = 1$. Then $D_k = \det A$, and we can write

$$\frac{1}{(\det A)^{1/(n-1)}} \le P \le \frac{1 + \sqrt{1 - (\det A)^{n-1}}}{1 - \sqrt{1 - (\det A)^{n-1}}}.$$
 (15)

This implies that

det $A \rightarrow 0$ if and only if $P \rightarrow \infty$,

det
$$A \rightarrow 1$$
 if and only if $P \rightarrow 1$.

Hence, the det A behaves essentially as the reciprocal of the P-condition number and may be said to constitute a reasonable condition number of its own for such matrices. This lends substance to the popular feeling that for a properly defined class of normalized matrices the smallness of its determinant is accompanied by difficulty in inversion.

A second observation relates to least square approximations. Express the general problem in the language of inner product spaces. We are required to solve $||y - \sum_{i=1}^{n} a_i x_i|| = \text{minimum}$, where y is given and x_i are independent elements. The normal equations have matrix $((x_i, x_j))$ where (p,q) designates the inner product. We can assume x_1 normalized: $(x_i, x_i) = ||x_i||^2 = 1$. The Gram matrix $((x_i, x_j))$ is positive definite symmetric and hence falls within the scope of our inequality with $k=1, c_1=n$. Hence if G is the Gram determinant $|(x_i, x_j)|$ we have

$$\frac{1}{G^{1/(n-1)}} \le P \le \frac{1 + \sqrt{1 - G}}{1 - \sqrt{1 - G}}.$$
(16)

The quantity G, which acts as a "measure" of linear independence of x_1, \ldots, x_n , therefore also serves as a condition number for the normal equations.

4. References

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